

Introduction to Data Structures

- Data Structures
- Stacks and Queues
- Linked Lists
- Dynamic Arrays

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- Dynamic Arrays

Data Structures

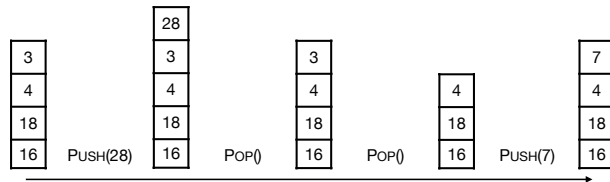
- **Data structure.** Method for organizing data for efficient access, searching, manipulation, etc.
- **Goal.**
 - Fast.
 - Compact
- **Terminology.**
 - **Abstract** vs. **concrete** data structure.
 - **Dynamic** vs. **static** data structure.

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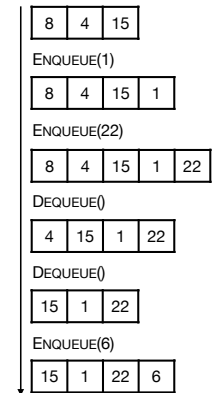
Stack

- **Stack.** Maintain dynamic sequence (stack) S supporting the following operations:
 - PUSH(x): add x to S.
 - POP(): remove and return the **most recently** added element in S.
 - ISEMPY(): return true if S is empty.



Queue

- **Queue.** Maintain dynamic sequence (queue) Q supporting the following operations:
 - ENQUEUE(x): add x to Q.
 - DEQUEUE(): remove and return the **earliest added** element in Q.
 - ISEMPY(): return true if Q is empty.

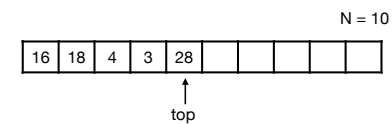


Applications

- **Stacks.**
 - Virtual machines
 - Parsing
 - Function calls
 - Backtracking
- **Queues.**
 - Scheduling processes
 - Buffering
 - Breadth-first searching

Stack Implementation

- **Stack.** Stack with **capacity** N
- **Data structure.**
 - Array S[0..N-1]
 - Index top. Initially top = -1
- **Operations.**
 - PUSH(x): Add x at S[top+1], top = top + 1
 - POP(): return S[top], top = top - 1
 - ISEMPY(): return true if top = -1.
 - Check for overflow and underflow in PUSH and POP.



Stack Implementation

• Time

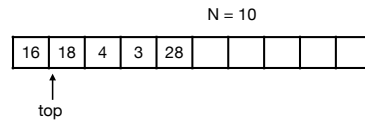
- PUSH in $O(1)$ time.
- POP in $O(1)$ time.
- ISEEMPTY in $O(1)$ time.

• Space.

- $O(N)$ space.

• Limitations.

- Capacity must be known.
- Wasting space.



Queue Implementation

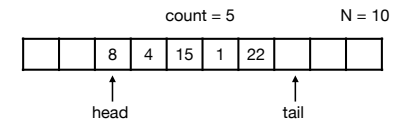
• Queue. Queue with capacity N.

• Data structure.

- Array $Q[0..N-1]$
- Indices head and tail and a counter.

• Operations.

- ENQUEUE(x): add x at $Q[\text{tail}]$, update count and tail *cyclically*.
- DEQUEUE(): return $Q[\text{head}]$, update count and head *cyclically*.
- ISEEMPTY(): return true if count = 0.
- Check for overflow and underflow in DEQUEUE and ENQUEUE.



Queue Implementation

• Time.

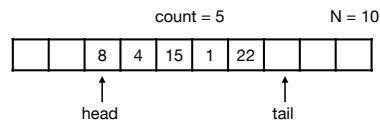
- ENQUEUE in $O(1)$ time.
- DEQUEUE in $O(1)$ time.
- ISEEMPTY in $O(1)$ time.

• Space.

- $O(N)$ space.

• Limitations.

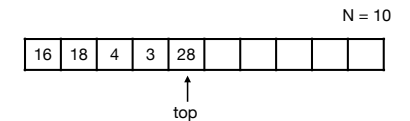
- Capacity must be known.
- Wasting space.



Stacks and Queues

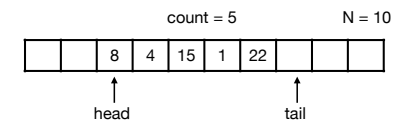
• Stack.

- Time. PUSH, POP, ISEEMPTY in $O(1)$ time.
- Space. $O(N)$



• Queue.

- Time. ENQUEUE, Dequeue, ISEEMPTY in $O(1)$ time.
- Space. $O(N)$



- Challenge. Can we get linear space and constant time?

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- **Linked Lists**
- Dynamic Arrays

Linked Lists

- **Linked lists.**
 - Data structure to maintain a **dynamic** sequence of items.
 - **Recursive** data structure. A linked list is either:
 - Empty
 - A reference to a **node** that has a reference to a linked list.
- **Node.**
 - An object that stores the item (or reference to the item) and the reference to a linked list.

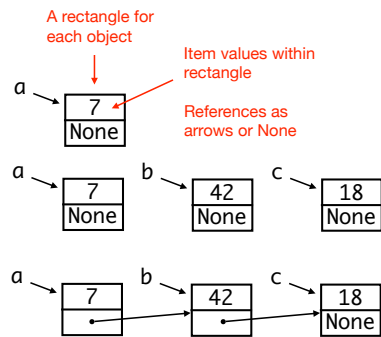
Linked Lists

```
class Node:  
    def __init__(self, item, next):  
        self.item = item  
        self.next = next
```

```
a = Node(7, None)
```

```
b = Node(42, None)  
c = Node(18, None)
```

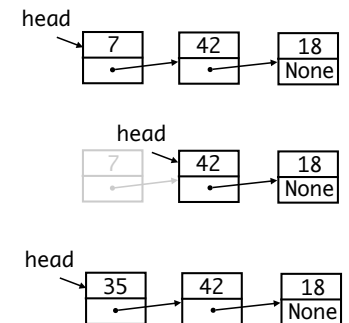
```
a.next = b  
b.next = c
```



Linked Lists

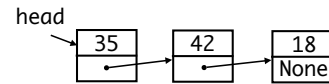
```
head = head.next
```

```
tmp = head  
head = Node(35, tmp)
```



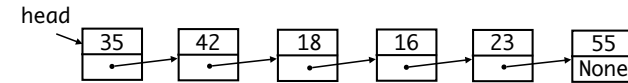
Linked Lists

```
x = head
while x != None:
    print(x.item)
    x = x.next
```



Linked Lists

- **Linked lists.**
 - Data structure to maintain a **dynamic** sequence of items.



- **Space.**
 - A list of n items uses $O(n)$ space.
- **Time.**
 - Insert at the head of the list in $O(1)$ time.
 - Delete at the head of the list in $O(1)$ time.
 - Traverse the list in $O(n)$ time.

Linked Lists

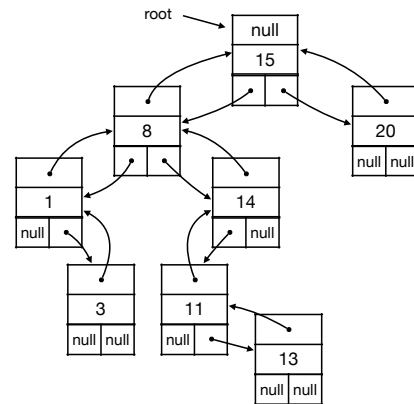
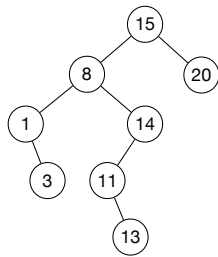
- **Exercise.** Consider how to implement stack and queue with linked lists efficiently.
- **Stack.** Maintain dynamic sequence (stack) S supporting the following operations:
 - **PUSH**(x): add x to S .
 - **POP**(): remove and return the **most recently** added element in S .
 - **ISEMPTY**(): return true if S is empty.
- **Queue.** Maintain dynamic sequence (queue) Q supporting the following operations:
 - **ENQUEUE**(x): add x to Q .
 - **DEQUEUE**(): remove and return the **earliest added** element in Q .
 - **ISEMPTY**(): return true if S is empty.

Linked Lists

- Stacks and queues using linked lists
- **Stack.**
 - **Time.** PUSH, POP, ISEMPTY in $O(1)$ time.
 - **Space.** $O(n)$
- **Queue.**
 - **Time.** ENQUEUE, DEQUEUE, ISEMPTY in $O(1)$ time.
 - **Space.** $O(n)$

Linked Lists

- **Linked list.** Flexible data structure to maintain sequence of elements.
- Other linked data structures: cyclic lists, trees, graphs, ...



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Stack Implementation with Array

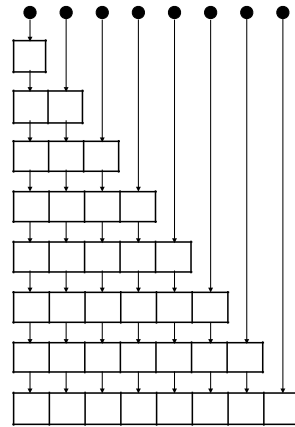
- **Challenge.** Can we implement a stack efficiently with arrays?
 - Do we need a fixed capacity?
 - Can we get linear space and constant time?

Dynamic Arrays

- **Goal.**
 - Implement a stack using arrays in $O(n)$ space for n elements.
 - As fast as possible.
 - Focus on PUSH. Ignore POP and ISEMPTY for now.
- **Solution 1**
 - Start with array of size 1.
- PUSH(x):
 - Allocate new array of size + 1.
 - Move all elements to new array.
 - Delete old array.

Dynamic Arrays

- **PUSH(x):**
 - Allocate new array of size + 1.
 - Move all elements to new array.
 - Delete old array.
- **Time.** Time for n PUSH operations?
 - ith PUSH takes $O(i)$ time.
 - \Rightarrow total time is $1 + 2 + 3 + 4 + \dots + n = O(n^2)$
- **Space.** $O(n)$
- **Challenge.** Can we do better?

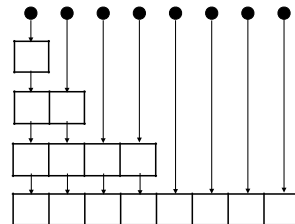


Dynamic Arrays

- **Idea.** Only copy elements some times
- **Solution 2.**
 - Start with array of size 1.
- **PUSH(x):**
 - If array is **full**:
 - Allocate new array of **twice the size**.
 - Move all elements to new array.
 - Delete old array.

Dynamic Arrays

- **PUSH(x):**
 - If array is **full**:
 - Allocate new array of **twice the size**.
 - Move all elements to new array.
 - Delete old array.
- **Time.** Time for n PUSH operations?
 - PUSH 2^k takes $O(2^k)$ time.
 - All other PUSH operations take $O(1)$ time.
 - \Rightarrow total time $< 1 + 2 + 4 + 8 + 16 + \dots + 2^{\log_2 n} + n = O(n)$
- **Space.** $O(n)$



Dynamic Arrays

- **Stack with dynamic array.**
 - n PUSH operations in $O(n)$ time and space.
 - Extends to n PUSH, POP and ISEMPTY operations in $O(n)$ time.
- Time is **amortized** $O(1)$ per operation.
- With more clever tricks we can **deamortize** to get $O(1)$ worst-case time per operation.
- **Queue with dynamic array.**
 - Similar results as stack.
- **Global rebuilding.**
 - Dynamic array is an example of **global rebuilding**.
 - Technique to make static data structures dynamic.

Stack and Queues

Data structure	PUSH	POP	ISEMPTY	Space
Array with capacity N	$O(1)$	$O(1)$	$O(1)$	$O(N)$
Linked List	$O(1)$	$O(1)$	$O(1)$	$O(n)$
Dynamic Array 1	$O(n)$	$O(1)^\dagger$	$O(1)$	$O(n)$
Dynamic Array 2	$O(1)^\dagger$	$O(1)^\dagger$	$O(1)$	$O(n)$
Dynamic Array 3	$O(1)$	$O(1)$	$O(1)$	$O(n)$

† = amortized

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