

Compressed Communication Complexity of Longest Common Prefixes

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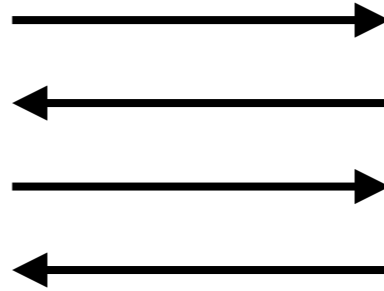
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Communication Complexity



$$f(A,B) = ?$$



A



B

- Bits of communication and number of rounds.
- Randomized.
- Public vs. private-coin.

Longest Common Prefixes

LCP(A,B)

A 01101....

B 01100....

$\ell = 4$

$|A| = |B| = n$

LCP^k(A,B)

A 01101....

B₁

B₂ 01100....

B₃

⋮

B_k

Equality

$A = B?$

A



P



0,1

B



P

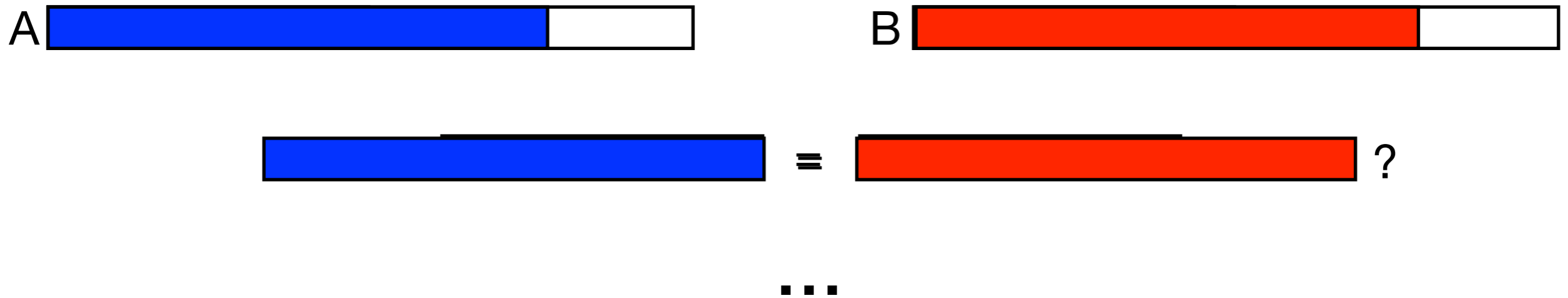


0,1

Theorem [Yao1979]. The public-coin randomized communication complexity of Equality is $O(1)$.

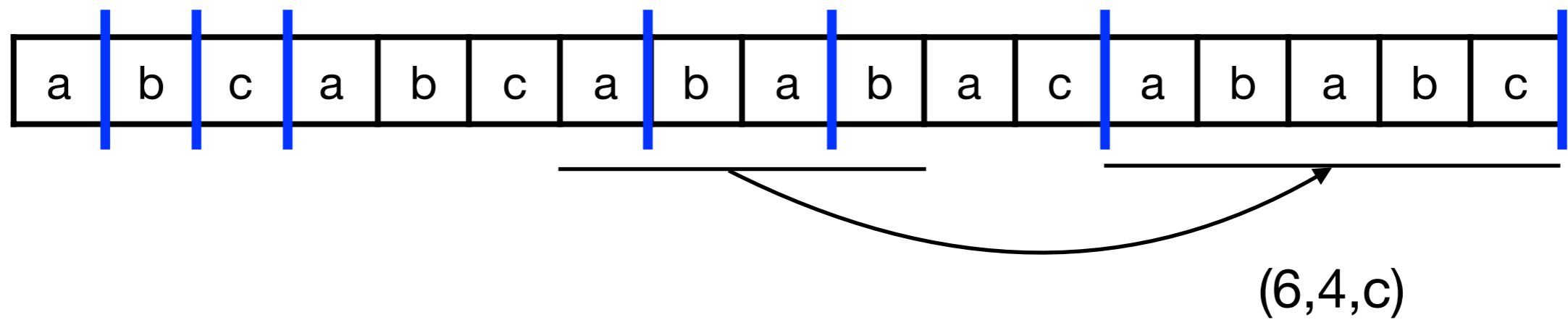
Longest Common Prefixes

LCP(A,B)



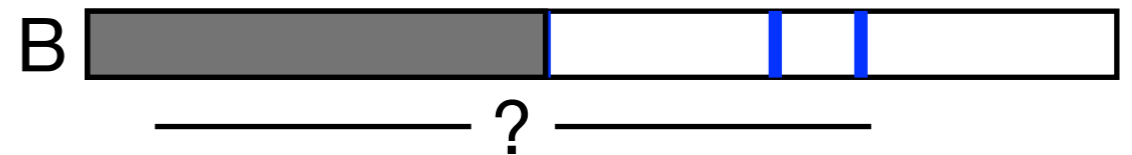
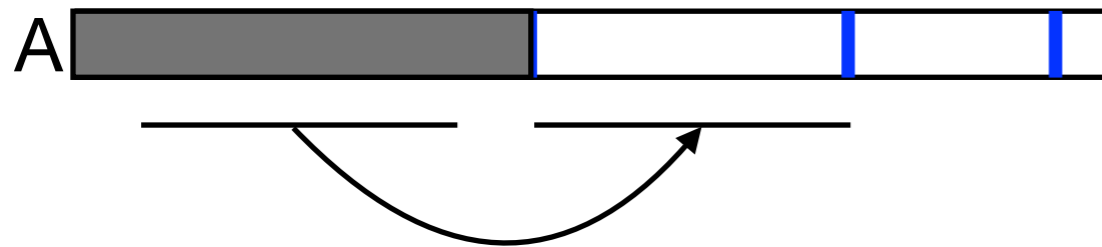
- $O(\log n)$ bits and $O(\log n)$ rounds. Optimal bound.
- Noisy binary search [FRPU1990].
- Noisy exponential search $\Rightarrow O(\log \ell)$ bits and $O(\log \ell)$ rounds.

Lempel-Ziv Compression



Compressed Longest Common Prefixes

LCP(A,B)



- $O(\log \ell)$ bits and $O(\log z)$ rounds.

	communication	rounds	ref
LCP	$O(\log n)$	$O(\log n)$	[N93]
	$O(\log \ell)$	$O(\log z)$	[new]
with selfref.	$O(\log \ell)$	$O(\log z + \log \log \ell)$	[new]
	$O(\log \ell + \log \log \log A)$	$O(\log z)$	[new]
LCP ^k	$O(\log z \log k + \log \ell)$	$O(\log z)$	[new]
	$O(\log z \log k + \log \ell)$	$O(\log z + \log \log \ell)$	[new]
with selfref.	$O(\log z \log k + \log \ell + \log \log \log A)$	$O(\log z)$	[new]