Domain Science & Engineering A Review of 10 Years Work and a Laudatio The ZCC Fest, 20 October 2017, Changsha, China

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Abstract. A personal account is given of my scientific work since I retired 10 years ago. This work centers around a new dimension to computing science: that of domain science & engineering. By a domain we shall understand a rationally describable segment of a human assisted reality, i.e., of the world, its physical parts, and living species. These are endurants ("still"), existing in space, as well as perdurants ("alive"), existing also in time. Emphasis is placed on "human-assistedness", that is, that there is at least one (man-made) artifact and that humans are a primary cause for change of endurant states as well as perdurant behaviours. Section 8 brings my laudatio.

1 Introduction

I survey recent work in the area of domain science & engineering¹.

A strict interpretation of the *triptych* of software engineering dogma suggests that software development "ideally" proceeds in three phases:

- First a phase of *domain engineering* in which an analysis of the application domain leads to a description of that domain.
- Then a phase of requirements engineering in which an analysis of the domain description leads to a prescription of requirements to software for that domain.
- And, finally, a phase of *software design* in which an analysis of the requirements prescription leads to software for that domain.

¹ It is appropriate, at this point, to state that my use of the term 'domain' is not related to that of *Domains and Processes* such as in the *Proceedings of 1st International Symposium on Domain Theory, Shanghai, China, October 1999, eds.: Klaus Keimel, Zhang Guo-Qiang, Liu Ying-Ming and Chen Yi-Chang. Springer Science + Business Media, New York, 2001.*

We see domain science & engineering as a discipline that need not be justified as a precursor to requirements engineering. Just as physicists study nature, irrespective of engineering, so we can study manifest domains irrespective of computing.

1.1 Recent Papers and Reports

Over the last decade I have iterated a number of investigations of aspects of this *triptych* dogma. This has resulted in a number of papers (and revised reports):

 Manifest Domains: Analysis & Description (2018, 2014) 	[29, 35]
 Domain Facets: Analysis & Description (2018. 2008) 	[31, 12]
 From Domains to Requirements (2018, 2008) 	[25, 8]
 Formal Models of Processes and Prompts (2014,2017) 	[23, 20]
 To Every Domain Mereology a CSP Expression (2017, 2009) 	[33, 10]
 A Philosophy of Domain Science & Engineering (2018) 	[30]

[30], a report, is the most recent.

1.2 Recent Experiments

Applications of the domain science and engineering outlined in [29]–[8] are exemplified in reports and papers on experimental domain analysis & description. Examples are:

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Urban Planning [41],
Documents [28],
Credit Cards [22],
Weather Information Systems [26],
The Tokyo Stock Exchange [34],
Pipelines [18],
Road Transportation [19],
Web/Transaction-based Software [14],
"The Market" [4],
Container [Shipping] Lines [7],
Railway Systems [3, 37, 5, 51, 56].
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1.3 My Emphasis on Software Systems

An emphasis in my work has been on research into and experiments with application areas that required seemingly large scale software. Not on tiny, beautiful, essential data structures and algorithms.

I first worked on the proper application of formal methods in software engineering at the IBM Vienna Laboratory in the early 1970s. That was to the formalisation of the semantics of IBMs leading programming language then, PL/I, and to a systematic development of a compiler for that language. The latter never transpired.

Instead I got the chance to formulate the stages of development of a compiler from a denotational semantics description to so-called "running code" [2, 1977]. That led, from 1978 onward, to two MSc students and a colleague and I working on a formal description of the *CCITT Communications High Level Language*, *CHILL* and its compiler [1, 46]. And that led, in 1980, to five MSc students of ours producing a formal description of a semantics for the *US DoD Ada programming language*, *Ada* [40]. And that led to the formation of *Dansk Datamatik Center* [38] which embarked on the CHILL and Ada compiler developments [42, 50]. To my knowledge that project which was on time, at budget, and with a history of less that 3% cost of original budget for subsequent error correction over the first 20 years of use of that compiler was a first, large, successful example of the systematic use of formal methods in large scale (42 man years) software development.

1.4 How Did We Get to Domain Science & Engineering?

So that is how we came from the semantics of programming languages to the semantics of human-centered, manifest application domain software development. Programming language semantics has to do with the meaning of abstract concepts such as programs, procedures, expressions, statements, GOTOs, labels, etc. Domain semantics, for manifest domains, in so far as we can narrate and formalize it, or them, must capture some "meanings" of the manifest objects that we can touch and see, of the actions we perform on them, and of the sentences by means of which we talk about those phenomena in the domain.

1.5 Preliminaries

We need formulate a few characterisations.

Method & Methodology: By a *method* I understand a set of principles for selecting and applying techniques and tools for constructing a manifest or an abstract artifact.

By methodology I understand the study and knowledge of methods.

My contributions over the years have contributed to methods for software design and, now, for the last many years, methods for domain analysis & description.

In my many experiments with domain analysis & description, cf. Sect. 5 on Page 22, I have found that I often let a so-called "streak of creativity" enter my analysis & description – and, as a result I get stuck in

my work. Then I recall, ah!, but there are these principles, techniques and tools for analysis & description, and once I apply them, "strictly", i.e., methodically, I am back on the track, and, in my view, a more beautiful description emerges!

Computer & Computing Sciences: By *computer science* I understand the study and knowledge about the things that can exist inside computing devices.

By computing science I understand the study and knowledge about how to construct the things that can exist inside computing devices. Computing science is also often referred to as programming methodology. My work is almost exclusively in the area of computing science.

A Triptych of Informatics: Before software can be designed we must have a firm grasp on its/their requirements. Before requirements can be prescribed we must have a firm grasp on their basis: the domain. We therefore see informatics as consisting of

- domain science & engineering,
- requirements science & engineering, and
- programming methodology.

This paper contributes to the establishment of domain science & engineering, while hinting that requirements science & engineering can benefit from the relation between the two [25,8]. How much of a domain must we analyse & describe before we attempt the second and third phases of the triptych? When this question is raised, after a talk of mine over the subject, and by a colleague researcher & scientist I usually reply: As large a domain as possible! This reply is often met by this comment (from the audience) Oh! No, that is not reasonable! To me that comment shows either or both of: the questioner was not asking as a researcher/scientist, but as an engineer. Yes, an engineer needs only analyse & describe up to and slightly beyond the "border" of the domain-of-interest for a current software development – but a researcher cum scientist is, of course, interested not only in a possible requirements engineering phase beyond domain engineering, but is also curious about the larger context of the domain, in possibly establishing a proper domain theory, etc.

1.6 The Papers

IM²HO I consider the first of the papers reviewed, [29], my most important paper. It was conceived of last², after publication of three of the other papers [12, 8, 16]. Experimental evidence then necessitated extensive revisions to these other papers, resulting in [31, 25, 24].

1.7 Structure of This Paper

Section 2 reviews [29, Analysis & Description Prompts], and Sect. 3 reviews related science and methodology papers. [31, Domain Facets] (Sect. 3.1), [25, From Domains to Requirements] (Sect. 3.2), [23, An Analysis & Description Process Model] (Sect. 3.3), and [33, From Mereologies to Lambda-Expressions] (Sect. 3.4), Finally, Sect. 4 briefly reviews [30, A Philosophy Basis] work-in-progress.

2 Manifest Domains: Analysis & Description [29]

This work grew out of many years of search for principles, techniques and tools for systematically analyzing and describing manifest domains. By a manifest domain we shall understand a domain whose entities we can observe and whose endurants we can touch!

2.1 A Domain Ontology

Parts, Components and Materials: The result became a calculus of analysis and description prompts³. These prompts are tools that the domain analyser & describer uses. The domain analyser & describer is in the domain, sees it, can touch it, and then applies the prompts, in some orderly fashion, to what is being observed. So, on one hand, there is the necessarily informal domain, and, on the other hand, there are the seemingly formal prompts and the "suggestions for something to be said", i.e., written down: narrated and formalised. Figure 1 on the next page suggests a number of analysis and description prompts. The domain analyser & describer is "positioned" at the top, the "root". If what is observed can be conceived and described then it is an entity. If it can be described as a "complete thing" at no matter which given snapshot of time then it is an endurant. If it is an entity but for which only a fragment exists if we look

² Publication [13, 15] is a predecessor of [35] which is then a predecessor of [29].

³ Prompt, as a verb: to move or induce to action; to occasion or incite; inspire; to assist (a person speaking) by "suggesting something to be said".

at or touch them at any given snapshot in time, then it is a perdurant. Endurants are either discrete or continuous. With discrete endurants we can choose to associate, or to not associate mereologies⁴. If we do we shall refer to them as parts, else we shall call them components. The continuous endurants we shall also refer to as (gaseous or liquid) materials. Parts are

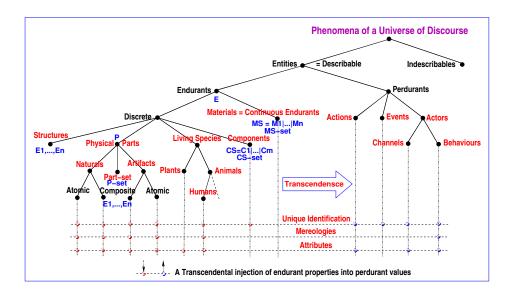


Fig. 1. Domain Ontology

either atomic or composite and all parts have unique identifiers, mereology and attributes. Atomic parts may have one or more components and/or one or more materials

If the observed part, p:P, is composite then we can observe the part sorts, $P_1, P_2, ..., P_m$ of p: observe_part_sorts(p) which yields the informal and formal description: Narrative: ... Formal: type $P_1, P_2, ..., P_m$, value obs_ P_i : $P \to P_i$, repeated for all m part sorts P_i s"!

Part sorts may have a concrete type: $has_concrete_type(p)$ in which case $observe_concrete_part_type(p)$ yields Narrative: ... Formal: type: T = P-set, value obs_T: $P \to K$ -set – where K-set is one of the concrete type forms, and where K is some sort.

Materials have types (i.e., are of sorts): M_i . Observing the (one) material, of type M, of an endurant e of sort E is expressed as *obs_materials*(e)

⁴ — 'mereology' will be explained next

which yields some narrative and some formal description text: Narrative: ... Formal: type M value obs_M: $E \to M$. The narrative text (...) narrates what the formal text expresses⁵.

Components, i.e., discrete endurants for whom we do not consider possible mereologies or attributes, can be observed from materials, m:M, or are just observed of discrete endurants, $e:E:obs_components(em)$ which yields the informal and formal description: Narrative: ... Formal: type: $C_1, C_2, ..., C_n$ value obs $_C_i$: $(E|M) \to C_i$ repeated for all n component sorts Cs" to the formal text!

• • •

The above is a pedagogic simplification. As shown in Fig. 1 on the facing page there are not only parts. There are also *living species: plants* and *animals*, including *humans*. And, because there are humans in the domains, parts and materials are either *natural* or *artifacts* (man-made). Humans create artifacts, usually with an *intent*. Humans have *intents*, and artifacts "possess" *intents*. Intents are like *attributes*, see below.

• • •

We have just summarised the analysis and description aspects of endurants in *extension* (their "form"). We now summarise the analysis and description aspects of endurants in *intension* (their "contents"). There are three kinds of intensional *qualities* associated with parts, two with components, and one with materials. Parts and components, by definition, have *unique identifiers*; parts have *mereologies*, and all endurants have *attributes*.

Unique identifiers: Unique identifiers are further undefined tokens that uniquely identify parts and components. The description language observer uid_P , when applied to parts p:P yields the unique identifier, $\pi:\Pi$, of p. The $observe_part_sorts(p)$ invocation yields the description text: ... [added to the narrative and] $type \Pi_1, \Pi_2, ..., \Pi_m$; value $uid_P : P_i \to \Pi_i$, repeated for all m part sorts P_i s and added to the formalisation.

Mereology: Mereology is the study and knowledge of parts and part relations. The mereology of a part is an expression over the unique identifiers

⁵ – not how it expresses it, as, here, in the RAISE [44] Specification Language, RSL [43].

of the (other) parts with which it is related, hence $mereo_P: P \to \mathcal{E}(\Pi_j, ..., \Pi_k)$, $\mathcal{E}(\Pi_j, ..., \Pi_k)$ is a type expression. The $observe_part_sorts(p)$ invocation yields the description text: ... [added to the narrative and] value $mereo_P_i: P_i \to \mathcal{E}_i(\Pi_{i_j}, ..., \Pi_{i_k})$ [added to the formalisation]

Example: The mereologies, (i, o), of pipe units in a pipeline system thus express, for each kind of pipe unit, whether it is a well, a linear pipe, a fork, a join, a pump, a valve, or a sink, the identities of the zero, one or two pipe units that it is "connected" to on the input, i, respectively the output, o, side: for well (0,1), for pipe (1,1), for fork (1,2), for join (2,1), for valve (1,1), for pump (1,1), for sink (1,0) units

Attributes: Attributes are the remaining qualities of endurants. The analysis prompt obs_attributes applied to an endurant yields a set of type names, $A_1, A_2, ..., A_t$, of attributes. They imply the additional description text: Narrative: ... Formal: type $A_1, A_2, ..., A_t$ value attr_A_i: $E \rightarrow A_i$ repeated for all t attribute sorts A_i s! Examples: Typical attributes of a person are Gender, Weight, Height, Birth date, etcetera. Dynamic and static attributes of a pipe unit include current flow into the unit, per input, if any, current flow out of the unit, per output, if any current leak from the unit, guaranteed maximum flow into the unit, guaranteed maximum flow into the unit, etcetera. Michael A. Jackson [49] categorizes attributes as either static or dynamic, with dynamic attributes being either inert, reactive or active. The latter are then either autonomous, biddable or programmable. This categorization has a strong bearing on how these (f.ex., part) attributes are dealt with when now interpreting parts as behaviours.

2.2 From Manifest Parts to Domain Behaviours

[35] then presents an interpretation, τ , which to manifest *parts* associate *behaviours*. These are then specified as CSP [48] *processes*. This interpretation amounts to a *transcendental deduction!*

The Transcendental Deduction Idea — by means of an example: The term train can have the following "meanings": The train, as an endurant, parked at the railway station platform, i.e., as a composite part. The train, as a perdurant, as it "speeds" down the railway track, i.e., as a behaviour. The train, as an attribute,

Atomic Parts: Atomic parts translate into their core behaviours: b_{core}^{patom} . The *core* behaviours are tail recursively defined, that is, are cyclic. $b_{core}^{patom}(...) \equiv (..., b_{core}^{patom}(...))$ where (...) indicate behaviour (i.e., function) arguments.

Composite Parts: A composite part, p, "translates", τ , into the parallel composition of a core behaviour: $b_{core}^{p_{comp}}(...)$, for part p, with the parallel composition of the translations, τ , for each of the parts, $p_1, p_2, ..., p_m$, of p, $(\tau(p_1)||\tau(p_2)||...||\tau(p_m))$ that is: $\tau(p)\equiv b_{core}^{p_{comp}}(...)||(\tau(p_1)||\tau(p_2)||...||\tau(p_m))$

Concrete Parts: The translation of concrete part set, t, types, $t: T = K - \mathbf{set}$, is $\tau(t) \equiv \|\{\tau(k_i) | k_i : K \cdot k_i \in t\}$.

Translation of Part Qualities (...): Part qualities, that is: unique identifiers, mereologies and attributes, are translated into behaviour arguments – of one kind or another, i.e., (...). Typically we can choose to index behaviour names, b by the unique identifier, id, of the part based on which they were translated, i.e., b_{id} . Mereology values are usually static, and can, as thus, be treated like we treat static attributes (see next), or can be set by their behaviour, and are then treated like we treat programmable attributes (see next), i.e., (...). Static attributes become behaviour definition (body) constant values. Inert, reactive and autonomous attributes become references to channels, say ch_dyn , such that when an inert, reactive and autonomous attribute value is required it is expressed as ch_dyn ?. Programmable and biddable attributes become arguments which are passed on to the tail-recursive invocations of the behaviour, and possibly updated as specified [with] in the body of the definition of the behaviour, i.e., (...).

2.3 Contributions of [29] – and Open Problems

For the first time we have, now, the beginnings of a calculus for developing domain descriptions. In [13, 15] we speculate on laws that these analysis & description prompts (i.e., their "meanings") must satisfy. With this calculus we can now systematically develop domain descriptions [41–56]. I am right now working on understanding issues of *implicit/explicit* semantics⁶ Since December 2017 I have revised [35] extensively: simplified it, extended it, clarified some issues, provided analysis & description techniques for channels and arguments, et cetera. The revised paper is [29]⁷.

⁶ Cf. http://impex2017.loria.fr/

 $^{^7}$ You can find it on the Internet: http://www.imm.dtu.dk/~dibj/2018/tosem/Bjorner-TOSEM.pdf.

3 Related Papers

3.1 Domain Facets: Analysis & Description [31, 12]

Overview By a domain facet we shall understand one amongst a finite set of generic ways of analyzing a domain: a view of the domain, such that the different facets cover conceptually different views, and such that these views together cover the domain.

[31] is an extensive revision of [12]. Both papers identify the following facets: intrinsics, support technologies, rules & regulations, scripts, license languages, management & organisation, and human behaviour. Recently I have "discovered" what might be classified as a domain facet: classes of attribute semantics: the diversity of attribute semantics resolving the issue of so-called implicit and explicit semantics. I shall not cover this issue in this talk.

Intrinsics: By domain intrinsics we shall understand those phenomena and concepts of a domain which are basic to any of the other facets, with such domain intrinsics initially covering at least one specific, hence named, stakeholder view.

Support Technology: By a domain support technology we shall understand ways and means of implementing certain observed phenomena or certain conceived concepts.

Rules and Regulations: By a domain rule we shall understand some text (in the domain) which prescribes how people or equipment are expected to behave when dispatching their duties, respectively when performing their functions.

By a domain regulation we shall understand some text (in the domain) which prescribes what remedial actions are to be taken when it is decided that a rule has not been followed according to its intention.

Scripts: By a domain script we shall understand the structured, almost, if not outright, formally expressed, wording of a procedure on how to proceed, one that possibly has legally binding power, that is, which may be contested in a court of law.

Management & Organisation: By domain management we shall understand such people (such decisions) (i) who (which) determine, formulate and thus set standards (cf. rules and regulations) concerning strategic, tactical and operational decisions; (ii) who ensure that these decisions are passed on to (lower) levels of management and to floor staff; (iii) who make sure that such orders, as they were, are indeed carried out; (iv) who handle undesirable deviations in the carrying out of these orders cum decisions; and (v) who "backstops" complaints from lower management levels and from "floor" staff.

By domain organisation we shall understand (vi) the structuring of management and non-management staff "overseeable" into clusters with "tight" and "meaningful" relations; (vii) the allocation of strategic, tactical and operational concerns to within management and non-management staff clusters; and hence (viii) the "lines of command": who does what, and who reports to whom, administratively and functionally.

Human Behaviour: By domain human behaviour we shall understand any of a quality spectrum of carrying out assigned work: (i) from careful, diligent and accurate, via (ii) sloppy dispatch, and (iii) delinquent work, (iv) to outright criminal pursuit.

Contributions of [31, 12] – and Open Problems: [31] now covers techniques and tools for analyzing domains into these facets and for their modeling. The issue of *license languages* are particularly intriguing. The delineations between the listed⁸ facets is necessarily not as precise as one would wish: we are dealing with an imprecise world, that of (manifest) domains. License languages are treated in [31].

3.2 From Domains to Requirements [25, 8]

Overview: [25] outlines a calculus of refinements and extensions which applied to domain descriptions yield requirements prescriptions. As for [35] the calculus is to be deployed by human users, i.e., requirements engineers. Requirements are for a *machine*, that is, the hardware and software to be developed from the requirements. A distinction is made between *domain*, *interface* and *machine requirements*. I shall briefly cover these in another order.

⁸ We have omitted a facet: *license languages*.

Machine requirements: Machine requirements are such which can be expressed using only technical terms of the machine: performance and dependability accessibility, availability, integrity, reliability, safety, security and robustness). and development requirements development process, maintenance, platform, management and documentation). Within maintenance requirements there are adaptive, corrective, perfective, preventive, and extensional requirements. Within platform requirements there are development, execution, maintenance, and demonstration requirements. Etcetera. [25] does not cover these. See instead [6, Sect. 19.6].

Domain Requirements: Domain requirements are such which can be expressed using only technical terms of the domain. The are the following domain-to-requirements specification transformations: projection, instantiation, determination, extension and fitting. I consider my work on these domain requirements issues the most interesting.

- 1: Projection: By a domain projection we mean a subset of the domain description, one which projects out all those endurants: parts, materials and components, as well as perdurants: actions, events and behaviours that the stake-holders do not wish represented or relied upon by the machine.
- 2: Instantiation: By domain instantiation we mean a refinement of the partial domain requirements prescription (resulting from the projection step) in which the refinements aim at rendering the endurants: parts, materials and components, as well as the perdurants: actions, events and behaviours of the domain requirements prescription more concrete, more specific.
- 3: Determination: By domain determination we mean a refinement of the partial domain requirements prescription, resulting from the instantiation step, in which the refinements aim at rendering the endurants: parts, materials and components, as well as the perdurants: functions, events and behaviours of the partial domain requirements prescription less non-determinate, more determinate.
- 4: Extension: By domain extension we understand the introduction of endurants and perdurants that were not feasible in the original domain, but for which, with computing and communication, and with new, emerging technologies, for example, sensors, actuators and satellites, there is the possibility of feasible implementations, hence the requirements, that what is introduced becomes part of the unfolding requirements prescription.

5: Fitting: Often a domain being described "fits" onto, is "adjacent" to, "interacts" in some areas with, another domain: transportation with logistics, health-care with insurance, banking with securities trading and/or insurance, and so on. The issue of requirements fitting arises when two or more software development projects are based on what appears to be the same domain. The problem then is to harmonize the two or more software development projects by harmonizing, if not too late, their requirements developments.

Interface Requirements: Interface requirements are such which can be expressed only by using technical terms of both the domain and the machine. Thus interface requirements are about that which is *shared* between the domain and the machine: *endurants* that are represented in machine storage as well as co-existing in the domain; *actions* and *behaviours* that are performed while interacting with phenomena in the domain; etc.

Contributions of [25, 8]: [25] does not follow the "standard division" of requirements engineering into systems and user requirements etcetera. Instead [25] builds on domain descriptions and eventually gives a rather different "division of requirements engineering labour" — manifested in the domain, the interface and the machine requirements paradigms, and these further into sub-paradigms, to wit: projection, instantiation, determination, extension and fitting. Some readers have objected to my use of the term refinement for the domain-to-requirements transformations.

3.3 Formal Models of Processes and Prompts [23, 20]

Overview: [35] outlines a calculus of prompts, to be deployed by human users, i.e., the domain analyzers & describers. That calculus builds on the assumption that the domain engineers build, in their mind, i.e., conceptually, a syntactical structure of the domain description, although, what the domain engineers can "see & touch" are semantic objects. A formal model of the analysis and description prompt process and of the meanings of the prompts therefore is split into a model for the process and a model of the syntactic and semantics structures.

A Summary of Analysis and Description Prompts

The Analysis Prompts:

[a]	is_entity	[f]	is_part	[k]	has_concrete_type
[b]	is_endurant	[g]	is_component	[1]	has_mereology
[c]	is_perdurant	[h]	is_material	[m]	has_components
[d]	is_discrete	[i]	is_atomic	[n]	has_material
[e]	is_continuous	[j]	is_composite	[0]	has_parts

The Description Prompts:

```
[1] observe_part_sorts [5] observe_attributes
[2] observe_concrete_type [6] observe_component_sorts
[3] observe_unique_identifier [7] observe_part_material_sort
[4] observe_mereology [8] observe_material_part_sorts
```

A Glimpse of the Process Model

Process "Management": Domain description involves the "generation" and use of an indefinite number of type (sort) names, Nm. The global, assignable variables αps and νps serve to hold the names of the sorts to be analysed, respectively the names of the sorts for which unique identifiers, mereologies and attributes have to be analysed and described.

```
type \mathsf{Nm} = \mathsf{PNm} \mid \mathsf{MNm} \mid \mathsf{KNm} variable \alpha \mathsf{ps} := [\Delta \mathsf{nm}] \ \mathsf{type} \ \mathsf{Nm-set} \nu \mathsf{ps} := [\Delta \mathsf{nm}] \ \mathsf{type} \ \mathsf{Nm-set} value \mathsf{sel\_and\_remove\_Nm} \colon \mathbf{Unit} \to \mathsf{Nm} \mathsf{sel\_and\_remove\_Nm}() \equiv \\ \mathsf{let} \ \mathsf{nm} : \mathsf{Nm} \bullet \mathsf{nm} \in \nu \mathsf{ps} \ \mathsf{in} \nu \mathsf{ps} := \nu \mathsf{ps} \setminus \{\mathsf{nm}\} \ ; \ \mathsf{nm} \ \mathsf{end}; \ \mathsf{pre} : \nu \mathsf{ps} \neq \{\}
```

Some Process Functions: The analyse_and_describe_endurants function is the major function. It invokes a number of other analysis & description functions. We illustrate two:

value

```
analyse_and_describe_endurants: \mathbf{Unit} \to \mathbf{Unit} analyse_and_describe_endurants() \equiv while \simis_empty(\nups) do
let nm = sel_and_remove_Nm() in
analyse_and_describe_endurant_sort(nm,\iota:nm) end end;
for all nm:PNm • nm \in \alphaps do if has_mereology(nm,\iota:nm)
then observe_mereology(nm,\iota:nm) end end
```

for all nm:Nm • nm $\in \alpha$ ps do *observe_attributes*(nm, ι :nm) end

```
analyse_and_describe_endurant_sort: NmVAL \rightarrow Unit analyse_and_describe_endurant_sort(nm,val) \equiv is_part(nm,val) \rightarrow analyse_and_describe_part_sorts(nm,val), is_material(nm,val) \rightarrow observe_material_part_sort(nm,val), is_component(nm,val) \rightarrow observe_component_sort(nm,val)
```

A Glimpse of the Syntax and Semantics Models We suggest a syntax and a semantics of domain descriptions.

The Syntactical Structure of Domains: First the syntax of domains – divided into the syntax of endurants parts, materials and components.

```
TypDef = PTypes \cup MTypes \cup KTypes
         PTypes = PNm \rightarrow PaTyp
         MTypes = MNm \rightarrow MaTyp
         \mathsf{KTypes} = \mathsf{KNm} \implies \mathsf{KoTyp}
       ENDType = PaTyp | MaTyp | KoTyp
          PaTyp == AtPaTyp \mid AbsCoPaTyp \mid ConCoPaTyp
        AtPaTyp :: mkAtPaTyp(s_qs:PQ,s_omkn:({|"nil"|}|MNn|KNm))
    AbsCoPaTyp :: mkAbsCoPaTyp(s_qs:PQ,s_pns:PNm-set)
                     \forall mkAbsCoPaTyp(pq,pns):AbsCoPaTyp • pns \neq {}
         axiom
    ConCoPaTyp :: mkConCoPaTyp(s_qs:PQ,s_p:PNm)
         MaTyp :: mkMaTyp(s\_qs:MQ,s\_opn:({|"nil"|}|PNm))
          KoTyp :: mkKoTyp(s_qs:KQ)
Then the syntax of the internal qualities of endurants:
            PQ = s_ui:UI \times s_me:ME \times s_atrs:ATRS
           ME == "nil" | mkUl(s_ui:Ul) | mkUlset(s_uil:Ul) | ...
         ATRS = ANm \rightarrow ATyp
    ANm, ATyp
           MQ = s_atrs:ATRS
            KQ = s\_uid:UI \times s\_atrs:ATRS
```

The Semantical Values of Domains: Corresponding, homomorphically, to these syntaxes are their semantics types:

```
ENDVAL = PVAL \mid MVAL \mid KVAL
      PVAL == AtPaVAL|AbsCoPVAL|ConCoPVAL
  AtPaVAL :: mkAtPaVAL(s_qval:PQVAL,
               s_omkvals:({|"nil"|}|MVAL|KVAL-set))
AbsCoPVAL :: mkAbsCoPaVAL(s\_qval:PQVAL,s\_pvals:(PNm \xrightarrow{m} PVAL))
    axiom ∀ mkAbsCoPaVAL(pqs,ppm):AbsCoPVAL•ppm≠[]
ConCoPVAL :: mkConCoPaVAL(s_qval:PQVAL,s_pvals:PVAL-set)
     MVAL :: mkMaVAL(s_qval:MQVAL,s_pvals:PVAL-set)
     KVAL :: mkKoVAL(s_qval:KQVAL)
Qualities: Semantic Types
    PQVAL = UIVAL \times MEVAL \times ATTRVALS
     UIVAL
    MEVAL == mkUIVAL(s_ui:UIVAL)|mkUIVALset(s_uis:UIVAL-set)|...
\mathsf{ATTRVALS} = \mathsf{ANm} \underset{\overrightarrow{m}}{\rightarrow} \mathsf{AVAL}
ANm, AVAL
    MQVAL = ATTRVALS
    KQVAL = UIVAL \times ATTRVALS
```

From Syntax to Semantics and "Back Again!" We define mappings from sort names to the possibly infinite set of values of the named type, and from endurant values to the names of their sort.

```
type
     Nm_to_ENDVALS =
          (\mathsf{PNm} \underset{\overrightarrow{m}}{\mathsf{PVAL}}.\mathsf{set}) \cup (\mathsf{MNm} \underset{\overrightarrow{m}}{\mathsf{MVAL}}.\mathsf{set}) \cup (\mathsf{KNm} \underset{\overrightarrow{m}}{\mathsf{KVAL}}.\mathsf{set})
     ENDVAL_to_Nm =
          (PVAL \xrightarrow{m} PNm) \cup (MVAL \xrightarrow{m} MNm) \cup (KVAL \xrightarrow{m} KNm)
value
     typval: TypDef \stackrel{\sim}{\to} Nm_to_ENDVALS
     typval(td) \equiv let \rho =
          [n\mapsto M(td(n))(\rho)|n:(PNm|MNm|KNm)\bullet n \in \mathbf{dom}\ td]\ \mathbf{in}\ \rho\ \mathbf{end}
     valtyp: Nm_to_ENDVALS \stackrel{\sim}{\to} ENDVAL_to_Nm
     valtyp(\rho) \equiv
          [v \mapsto n \mid n: (PNm \mid MNm \mid CNm), v: (PVAL \mid MVAL \mid KVAL) \bullet
               n \in \mathbf{dom} \ \rho \land v \in \rho(n)
     M: (PaTyp \rightarrow ENV \xrightarrow{\sim} PVAL-set)
          (MaTyp \rightarrow ENV \xrightarrow{\sim} MVAL-set)
          (KoTyp \rightarrow ENV \xrightarrow{\sim} KVAL-set)
```

The environment, ρ , of typval is the least fix point of the recursive equation. The crucial function is M, in the definition of typval. Examples of its definition, by part category, is given below.

```
value
```

```
\iota nm:Nm \equiv iota(nm)
iota: Nm \rightarrow TypDef \rightarrow VAL
iota(nm)(td) \equiv let val:(PVAL|MVAL|KVAL)•val\in(typval(td))(nm)
in val end
```

Analysis Functions: We exemplify the semantics functions for three analysis prompts.

value

```
is_endurant: Nm \times VAL \rightarrow TypDef \stackrel{\sim}{\rightarrow} \mathbf{Bool}
is_endurant(_,val)(td) \equiv val \in \mathbf{dom} \ valtyp(typval(td));
\mathbf{pre}: VAL is any value type
is_discrete: NmVAL \rightarrow TypDef \stackrel{\sim}{\rightarrow} \mathbf{Bool}
is_discrete(_,val)(td) \equiv (is\_PaTyp|is\_CoTyp)(td((valtyp(typval(td)))(val)))
is_part: NmVAL \rightarrow TypDef \stackrel{\sim}{\rightarrow} \mathbf{Bool}
is_part(_,val)(td) \equiv is\_PaTyp(td((valtyp(typval(td)))(val)))
```

Description Functions: We exemplify the semantics of one of the description prompts. The generated description $\mathbf{RSL\text{-}text}$ is enclosed within [" ... "].

```
variable
```

```
\begin{split} \tau := [\,] \ \mathbf{Text\text{-}set} \\ \mathbf{value} \\ & \text{observe\_part\_sorts: Nm} \times \mathsf{VAL} \to \mathsf{TypDef} \to \mathbf{Unit} \\ & \text{observe\_part\_sorts(nm,val)(td)} \equiv \\ & \text{let mkAbsCoPaTyp}(\_, \{\mathsf{P}_1, \mathsf{P}_2, ..., \mathsf{P}_n\}) \\ & = \mathsf{td}((\mathsf{valtyp}(\mathsf{typval}(\mathsf{td})))(\mathsf{val})) \ \mathbf{in} \\ & \tau := \tau \oplus [\, \text{"type } P_1, P_2, ..., P_n; \\ & \mathbf{value} \\ & \mathbf{obs\_part\_}P_1: \ nm \to P_1 \\ & \mathbf{obs\_part\_}P_2: \ nm \to P_2 \end{split}
```

```
proof obligation
                                         D: " ]
              \parallel \nu \mathit{ps} := \nu \mathit{ps} \oplus (["\ P_1, P_2, ..., P_n"] \setminus \alpha \mathit{ps})
              \parallel \alpha \textit{ps} := \alpha \textit{ps} \oplus [" P_1, P_2, ..., P_n"]
               end
       pre: is_AbsCoPaTyp(td((valtyp(typval(td)))(val)))
The M Function
 1 The meaning of an atomic part type expression,
      mkAtPaTyp((ui,me,attrs),omkn) in
     - \text{ mkAtPaTyp}(s\_qs:PQ,s\_omkn:(\{|"nil"|\}|MNn|KNm)),
     - is the set of all atomic part values,
         mkAtPaVAL((uiv,mev,attrvals),omkval) in
     - mkAtPaVAL(s_qval:(UIVAL\timesMEVAL\times(ANm \overrightarrow{m} AVAL)),
                        s_{omkvals:}({|"nil"|}|MVAL|KVAL-set)).
      a uiv is a value in UIVAL of type ui,
      b mev is a value in MEVAL of type me,
      c attrvals is a value in (ANm \overrightarrow{m} AVAL) of type (ANm \overrightarrow{m} ATyp), and
      d omkvals is a value in (\{|"nil"|\}|MVAL|KVAL-set):
           i either ''nil'',
          ii or one material value of type MNm,
         iii or a possibly empty set of component values, each of type KNm.
1. M: mkAtPaTyp((UI\timesME\times(ANm \overrightarrow{m}ATyp))\times({|"nil"|}|MVAL|KVAL-set))
             \rightarrowENV\stackrel{\sim}{\rightarrow}PVAL-\mathbf{set}
1.
   M(mkAtPaTyp((ui,me,attrs),omkn))(\rho) \equiv
1.
        { mkATPaVAL((uiv,mev,attrval),omkvals) |
1.
            uiv:UIVAL•type_of(uiv)=ui,
1a.
1b.
            mev:MEVAL•type_of(mev)=me,
            attrval:(ANm \overrightarrow{m} AVAL)•type_of(attrval)=attrs,
1c.
            omkvals: case omkn of
1d.
                  "nil" \rightarrow "nil",
1(d)i.
                   mkMNn(\underline{\hspace{0.1cm}}) \rightarrow mval:MVAL \bullet type\_of(mval) = omkn,
1(d)ii.
1(d)iii.
                   \mathsf{mkKNm}(\underline{\hspace{0.1cm}}) \to
                       kvals:KVAL-set \cdot kvals \subseteq \{kv|kv:KVAL \cdot type\_of(kval) = omkn\}
1(d)iii.
1d.
            end }
```

obs_part_ P_n : $nm \rightarrow P_n$;

Formula terms 1a–1(d)iii express that any applicable uiv is combined with any applicable mev is combined with any applicable attrval is combined with any applicable omkvals.

- 2 The meaning of an abstract composite part type expression,
 - mkAbsCoPaTyp((ui,me,attrs),pns) in
 - mkAbsCoPaTyp(s_qs:PQ,s_pns:PNm-set), is the set of all abstract, composite part values,
 - mkAbsCoPaVAL((uiv,mev,attrvals),pvals) in
 - mkAbsCoPaVAL(s_qval:(UIVAL \times MEVAL \times (ANm \overrightarrow{m} AVAL)), s_pvals:(PNm \overrightarrow{m} PVAL)).
 - a uiv is a value in UIVAL of type ui: UI,
 - b mev is a value in MEVAL of type me: ME,
 - c attrvals is a value in (ANm \overrightarrow{m} AVAL) of type (ANm \overrightarrow{m} ATyp), and
 - d pvals is a map of part values in ($PNm \xrightarrow{m} PVAL$), one for each name, pn:PNm, in pns such that these part values are of the type defined for pn.

```
M: mkAbsCoPaTyp((UI\timesME\times(ANm \rightarrow ATyp)),PNm-set)

ightarrow ENV \stackrel{\sim}{
ightarrow} PVAL-set
   M(mkAbsCoPaTyp((ui,me,attrs),pns))(\rho) \equiv
2.
2.
       { mkAbsCoPaVAL((uiv,mev,attrvals),pvals) |
2a.
           uiv:UIVAL•type_of(uiv)=ui
2b.
           mev:MEVAL•type_of(mev)=me,
2c.
           attrvals:(ANm \overrightarrow{m} ATyp)•type_of(attrsval)=attrs,
           pvals:(PNm \overrightarrow{m} PVAL) •
2d.
2d.
               pvals \in \{[pn \mapsto pval|pn:PNm,pval:PVAL \cdot pn \in pns \land pval \in \rho(pn)]\} \}
```

Contributions of [23] The contributions of [23] are to suggest and carry through a "formalisation" of the *conceptual*, *syntactical* and *semantical* structures *perceived* by the domain engineer, to formalise the *meaning of the informal* analysis & description prompts, and to formalise the possible sets of sequences of valid prompts.

3.4 To Every Manifest Domain Mereology a CSP Expression [33]

Overview In [35] we have shown how parts can be endowed with mereologies. Mereology, as was mentioned earlier, is the study and knowledge of "part-hood": of how parts are related parts to parts, and parts to "a whole". Mereology, as treated by us, originated with the Polish mathematician/logician/philosopher Stanislaw Lešhniewski.

An Axiom System for Mereology:

```
\begin{array}{ccc} \mathsf{part\_of:} & \mathbb{P}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \mathsf{proper\_part\_of:} & \mathbb{PP}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \mathsf{overlap:} & \mathbb{O}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \mathsf{underlap:} & \mathbb{U}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \mathsf{over\_crossing:} & \mathbb{OX}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \mathsf{under\_crossing:} & \mathbb{UX}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \mathsf{under\_crossing:} & \mathbb{UX}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \mathsf{proper\_overlap:} & \mathbb{PO}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \mathsf{proper\_underlap:} & \mathbb{PU}: \mathcal{P} \times \mathcal{P} \to \mathbf{Bool} \\ \end{array}
```

Let \mathbb{P} denote *part-hood*; p_x is part of p_y , is then expressed as $\mathbb{P}(p_x, p_y)$. (1) Part p_x is part of itself (reflexivity). (2) If a part p_x is part of p_y and, vice versa, part p_y is part of p_x , then $p_x = p_y$ (anti-symmetry). (3) If a part p_x is part of p_y and part p_y is part of p_z , then p_x is part of p_z (transitivity).

$$\forall p_x : \mathcal{P} \bullet \mathbb{P}(p_x, p_x) \tag{1}$$

$$\forall p_x, p_y : \mathcal{P} \bullet (\mathbb{P}(p_x, p_y) \land \mathbb{P}(p_y, p_x)) \Rightarrow p_x = p_y$$
 (2)

$$\forall p_x, p_y, p_z : \mathcal{P} \bullet (\mathbb{P}(p_x, p_y) \land \mathbb{P}(p_y, p_z)) \Rightarrow \mathbb{P}(p_z, p_z)$$
(3)

We exemplify one of the mereology propositions: *proper underlap*, \mathbb{PU} : p_x and p_y are said to properly underlap if p_x and p_y under-cross and p_y and p_x under-cross.

$$\mathbb{PU}(p_x, p_y) \stackrel{\triangle}{=} \mathbb{UX}(p_x, p_y) \wedge \mathbb{UX}(p_y, p_x)$$
(4)

A Model for the Axioms [33] now gives a model for parts: atomic and composite, commensurate with [35] and [23], and their unique identifiers, mereology and attributes and show that the model satisfies the axioms.

Contributions of [33] [33] thus contributes to a domain science, helping to secure a firm foundation for domain engineering.

4 Domain Science & Engineering: A Philosophy Basis [30]

My most recent work is documented in [30]. It examines the question:

— What must inescapably be in any domain description?

Another formulation is:

⁹ Our notation now is not RSL but a conventional first-order predicate logic notation.

 Which are the necessary characteristics of each and every possible world and our situation in it.

Recent works by the Danish philosopher Kai Sørlander [52–55] appears to direct us towards an answer.

Here is how it is done, in brief. On the basis of *possibility of truth*¹⁰ Sørlander establishes the logical connectors and from them the existence of a world with symmetry, asymmetry and transitivity. By a transcendental deduction Sørlander then reasons that *space* and *time*, inescapably, are "in the world"¹¹. Further logical reasoning and transcendental deductions establishes the inescapability of *Newton's 1st, 2nd and 3rd Laws*. And from that *kinematics*, *dynamics*, and *gravitational pull*. And so forth. Thus the worlds that can possibly be described must all satisfy the *laws of physics*.

This line of reasoning and deduction thus justifies the focus, in our calculi, on natural parts, components and materials.

But Sørlander goes on and reasons and transcendentally deduce the inescapable existence of *living species: plants* and *animals*, and, among the latter, *humans*. Because of reasoned characteristics of humans we inescapably have *artifacts: man-made parts components* and *materials*. Humans construct artifacts with an *intent*, an attribute of both humans and artifacts. These *shared intents* lead to a notion of *intentional "pull"* and so forth.

This line of reasoning and deduction thus justifies the inclusion, in our calculi, of living species and artifacts.

[30] is presently an approximately 90 page report. As such it is presently a repository for a number of "texts" related to the issue of "what must inescapably be in any domain description?" It may be expected that a far shorter paper may emerge.

¹⁰ Sørlander makes his *logical reasoning* and *transcendental deductions* on the basis of the *possibility of truth* – where Immanuel Kant [45], according to Sørlander, builds on the *possibility of self-awareness*, which is shown to lead to contradictions.

 $^{^{11}}$ Kant assumes space and time.

We shall here give an example of *intentional "pull"*: humans create automobiles and roads. An intention of automobiles is to drive on roads, and an intention of roads is to have automobiles move along roads. We can thus speak of the *traffic history of an automobile* as the time-stamped sequence of vehicle positions along roads, and of the *traffic history of a road* as the time-stamped sequence of vehicle positions along that road. Now, for the sum total of all automobiles and all roads the two consolidate histories must be identical. *It cannot be otherwise*.

5 The Experiments [41–56]

In order to test and tune the domain analysis & description method a great number of experiments were carried out. In our opinion, when applied to manifest domains, they justify the calculi reported in [35] and [23].

_	Urban Planning	[41],	- P	ipelines	[18],
_	A Space of Swarms of Drones	[27],	-R	oad Transportation	[19],
_	Documents	[28],	- V	Veb/Transaction-based Software	[14],
_	Credit Cards	[22],	_ "	The Market"	[4],
_	Weather Information Systems	[26],	- C	ontainer [Shipping] Lines	[7],
_	The Tokyo Stock Exchange	[34],	-R	Pailway Systems $[3, 37, 5, 51]$., 56].

6 Summary

We have identified a discipline of domain science and engineering. Its first "rendition" was applied to the semantics of programming languages and the development of their compilers [46, CHILL] and [42, Ada]. Domain science and engineering, as outlined here, is directed at a wider spectrum of "languages": the "meaning" of computer application domains and software for these applications. Where physicists model facets of the world emphasizing physical, dynamic phenomena in nature, primarily using differential calculi, domain scientists cum engineers emphasize logical and both discrete phenomena of man and human institutions primarily using discrete mathematics.

7 Acknowledgments

I am grateful to Prof. Zhan NaiJun for inviting me to the Zhou ChaoChen Fest and for inviting me to submit my talk as a paper for this Festschrift. I am grateful to my "old student", now Prof. Ji Wang for his arranging a wonderful stay in Changsha, my fourth visit to that great city, and for his fantastic cheerful welcome.

8 Laudatio

At the Zhou ChaoChen Fest dinner I gave a dinner speech. It is not about Zhou Chaochen's scientific life. But it is a laudatio expressed in admiration for a wonderful man and our lives together.

It was in 1981, in Beijing, 36 years ago. At the Institute of Computing Technology.
 On my first day of a three week visit. 30 lectures, 30 degrees Celsius. I liked it.

I was being received. All sat in soft cushioned armchairs along the walls.

I sat to the right of this wonderful man, Xu KongShi.

During our conversation I queried about a young researcher, Zhou ChaoChen.

Tony Hoare had told me to watch out for him.

So I did, with an invitation letter, right in my pocket, for him to visit my dept.

Asked Xu KongShi as to the whereabouts of Zhou?

And he smiled: right next to the right of you!

That became the first day of a 36 year acquaintance. Almost half of our life-times! Zhou came to visit us, 3 months every other winter. It was during the 1980s. What a wonderful time, for me, for my colleagues and for our students. One time I asked him to tutor a young MSc student. She performed brilliantly. It was something about "the meeting calendar problem". Even Zhou was impressed. Perhaps he has forgotten it now. When I took him to the airport, some weeks later. I told him that Ulla, that was her name, was a great granddaughter of Niels Bohr. Zhou appreciated then that I only told him then.

- For the 1989 visit I had "stipulated" that Zhou bring his family. Three months to Lyngby, three months to Oxford. And Zhou kindly agreed. All was

But a certain incident early that June caused us all concern.

Yet, on July 1st that year the whole family arrived.

Zhou wasn't keen to return to China.

I speed-dialled Tony's Oxford number.

"Tony on the line" was the reply

"Tony: Zhou is with me here, in my office in Lyngby."

"Hello Zhou"

"Hello Tony".

"Tony, I have just offered Zhou a three year appointment."

Well I hadn't, but there it was, and Zhou got listening.

"8 months a year here at Lyngby. 4 with you at Oxford."

Tony's reply: "Well, I had got it wrong, the other way around".

"Let Zhou decide", I replied, and Zhou said: "It is as Dines proposes."

Those became three great years, at Lyngby and at Oxford.

Zhang Yi Ping and children lodged in Oxford - Zhou commuting. Science progressing. It was at a ProCoS meeting in Viborg.

E.V.Sørensen had given a talk on signal transitions of electric circuits.

The concept of 'duration' was mentioned.

Afterwards I saw Zhou, A.P. and Tony, in an adjacent room.

Discussing, standing at the white board, scribbling.

And "The Duration Calculus" was born.

The following year I was asked to become Director of UNU-IIST.

On the flight home, in May 1991, from Japan, via a visit to Macau I decided to ask Zhou to join me in Macau.

And a year later, things take time in international affairs, we began.

With Zhou in charge of theory and I of engineering, an institute was built.

After my five years followed Zhou's five years.

Some of you, in this room, can look back at defining years at UNU-IIST.

I returned to Lyngby and eventually Zhou to Beijing.

- The Duration Calculus took root.
 - Painstakingly a theory was cemented and applications realized.
 - The ProCoS project and UNU-IIST played an important rôle in this.
 - But at the core of all this was Zhou ChaoChen.
- Dear Zhou:
 - Thank you for your tremendous contributions to science.
 - Thank you for inspiring generations of scientists.
 - Thank you for hosting our daughter, Charlotte, the fall of 1986 31 years ago!
 - Thank you for putting our son, Nikolaj, on the road to science also 31 years ago!

9 Bibliography

9.1 Bibliographical Notes

In the last ten years I have also worked on related topics:

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- Compositionality: Ontology and Mereology of Domains 13 , [36] 2008,
- Domain Science & Engineering, [13, 15] 2010,
- Computation for Humanity: Domain Science and Engineering, [17] 2012,
- 40 Years of Formal Methods Obstacles and Possibilities¹⁴, [39] 2014,
- Domain Engineering A Basis for Safety Critical Software, [21] 2014,
- Implicit and Explicit Semantics and the Domain Calculi, [32] 2017.

Work on these papers and on the many, extensive experiments has helped solidify the basic domain analysis & description method.

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 $[\]overline{}^{13}$ with Asger Eir

¹⁴ with Klaus Havelund

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