

# **Computational Tools for Data Science**

Week 6:

Clustering



- What is it: Task of grouping objects s.t. objects of the same group are more "similar" to each other than objects from different ones.

  The groups are called "clusters". In general they may overlap.
- For which data: numerical, usually a lot of data and high-dimensional

### Usually:

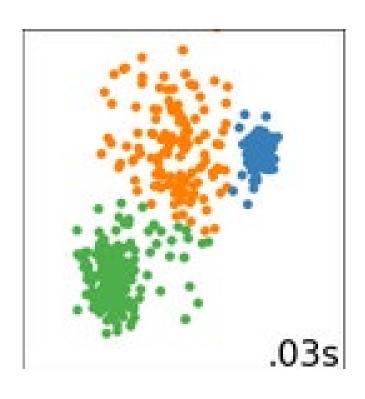
- Objects are points (or vectors) in a high-dimensional (vector) space.
- "Similarity" is defined via some distance measure, e.g.:
  - » vectors: Euclidean distance, Cosine
  - » Jaccard similarity (for sets), edit distance (strings)

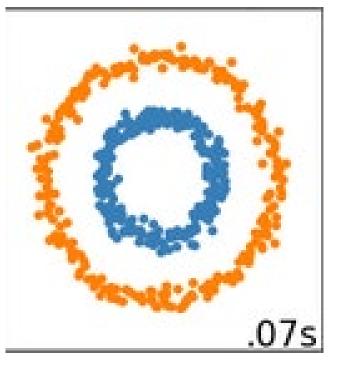


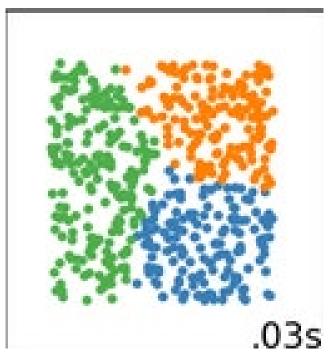
### Reminder: Distance

- **Definition (distance** d): Let V be a (top.) space. A map  $d: V \times V \to \mathbb{R}$  is called a *distance* if the following properties hold:
  - 1) For all  $u, v \in V$ :  $d(u, v) \ge 0$  and d(u, v) = 0 iff u = v.
  - 2) For all  $u, v \in V$ : d(u, v) = d(v, u)
  - 3) For all  $u, v, w \in V$ :  $d(u, v) + d(v, w) \ge d(u, w)$ . [Triangular inequality]











### **Examples:**

- In general: categorisation, classification and segmentation
  - » But also: finding anomalies, recognition
- Clustering stelar objects (clusters of stars)
- Group people by interest in social media (via #s, followers, topics)
- Group people / customer by consumed goods (e.g. music, movies, food)
- Find categories for goods by customers who bought them
- Identify topics for documents via shared words



**Example:** Categorise goods, say **movies**, via some set *C* of customers

- Represent each movie as |C| dimensional 0-1 vector, one dimension per customer from C.
- An entry / component in this vector is 1 iff the corresponding customer has watched / purchased the movie.
- > This yields thousands or even millions of dimensions (e.g. Amazon)
- Idea behind: Taste in movies causes us to buy / watch similar movies.



### **Hierarchical Clustering**

### A) Agglomerative (bottom up):

- 1. Initialise each data point as a (singleton) cluster.
- 2. Repeatedly combine two "most similar / closest" clusters to a new one.
- 3. Stop w.r.t. certain criterion (e.g. # of clusters, density of clusters, ...)



## **Hierarchical Clustering**

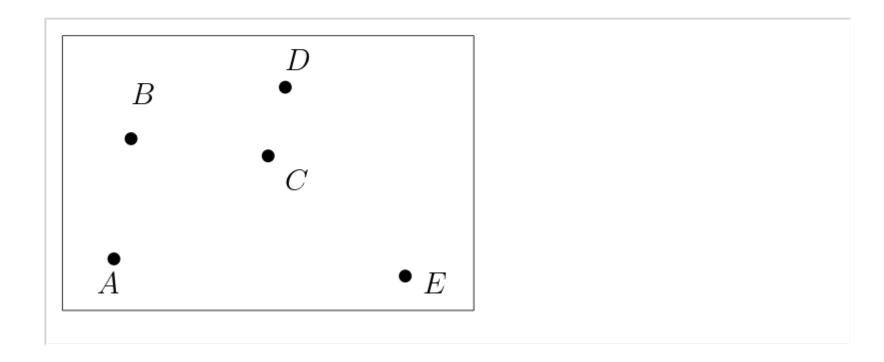
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### B) Divisive (top down):

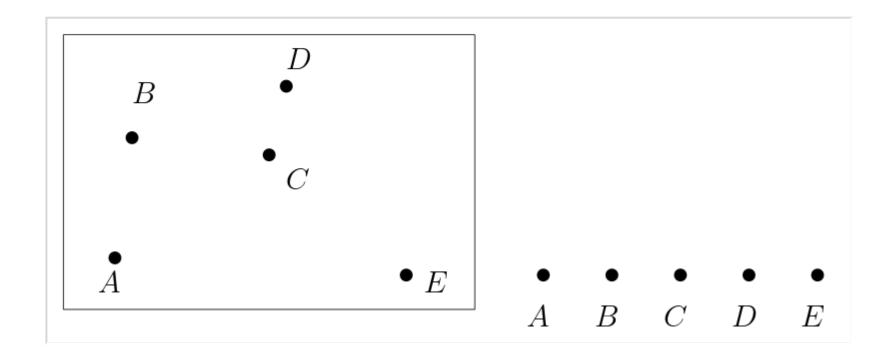
- 1. Initialise the whole data set as one big cluster.
- 2. Recursively split up cluster which is least "dense / connected".
- 3. Stop w.r.t. certain criterion (e.g. # of clusters, density of clusters, ...)





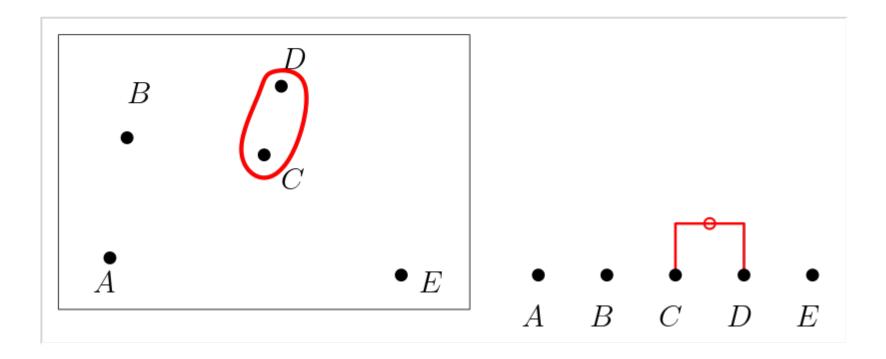
Resulting tree structure called dendrogram.





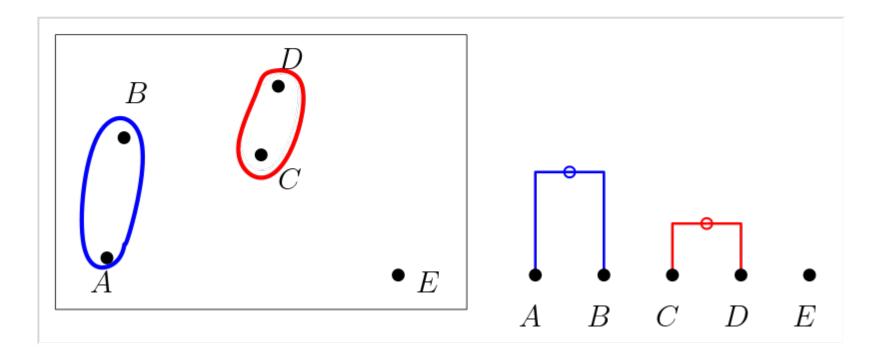
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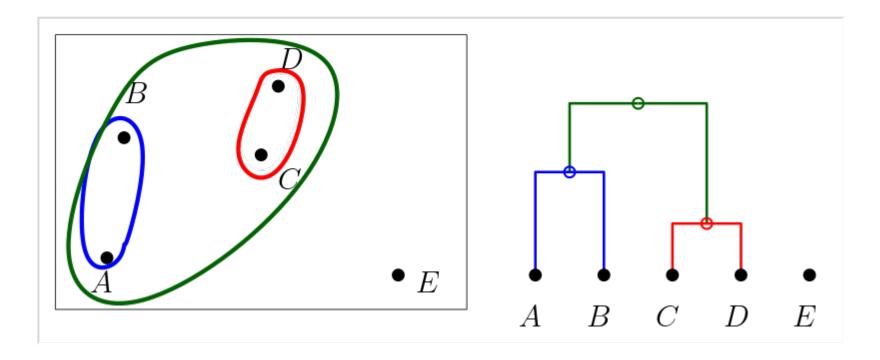


Resulting tree structure called dendrogram.

Use case: E.g. in phylogenetics with no huge data (phylogenetic tree)

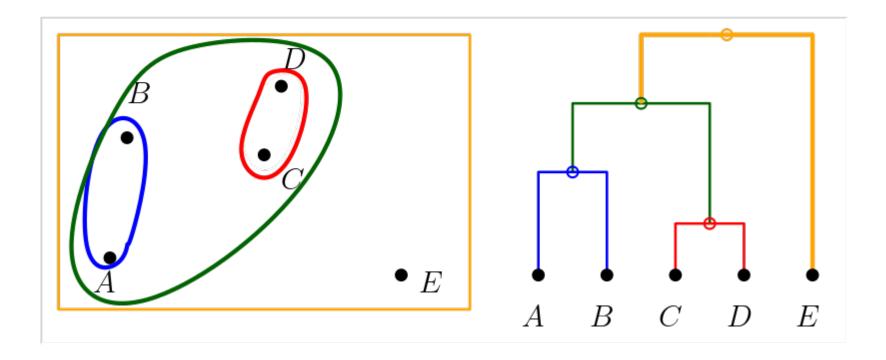
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Resulting tree structure called dendrogram.





Resulting tree structure called dendrogram.



## (Hierarchical) Clustering

### 3 important questions to answer:

- a) How do we represent clusters?
- b) How do we determine "similarity / closeness" of clusters?
- c) When do we stop the algorithm?



#### **Euclidean case:**

Let C be a cluster of n many d-dimensional data points, say in  $\mathbb{R}^d$ . We define the **centroid** cent(C) of C as follows:

$$cent(C) = \frac{1}{n} \sum_{c \in C} c$$

Think of **centroid** as a **centre of gravity**.



#### Non-Euclidean case:

Let C be a cluster of n many data points.

**Define** the clustroid clust(C) of C as a data point (from C) that is "closest" to the other points in C.

Some possible meanings of "closest":



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- Smallest sum of squared distances to other points in C.

$$\rightarrow \min_{c \in C} \sum_{v \in C} d^2(c, v)$$



Idea: Replace each cluster by its centroid (Euclidean) or clustroid (Non-Euclidean).

Measure distance of two clusters C and C' to each other by:

**Euclidean setting:** d(cent(C), cent(C'))

**Non-Euclidean setting:** d(clust(C), clust(C'))



Other ways to measure:

• Define distance between two clusters C and C' as the minimum distance between any data point from C and any data point in C'.

- Set a notion of "cohesion" (or connectivity, density, ...)
  - → Merge two cluster whose union is most cohesive.



Set a notion of "cohesion" (or connectivity, density, ...)

→ Merge two cluster whose union is most cohesive.

Possible notions of "cohesion":

A) Minimum diameter

For a cluster C define its diameter diam(C) as:

$$diam(C) = \max_{c,c' \in C} d(c,c')$$



Set a notion of "cohesion" (or connectivity, density, ...)

→ Merge two cluster whose union is most cohesive.

Possible notions of "cohesion":

B) Minimum average distance (between two points in the cluster)

Let C be a cluster with  $n \coloneqq |C|$  many points in it. Further let  $\binom{C}{2}$  denote the set of all pairs of points in C. Then the average distance avd(C) of C is:

$$avd(C) \coloneqq \frac{2}{n(n-1)} \sum_{\{c,c'\} \in \binom{C}{2}} d(c,c')$$



Set a notion of "cohesion" (or connectivity, density, ...)

→ Merge two cluster whose union is most cohesive.

Possible notions of "cohesion":

C) Minimum density

E.g. define the density of a cluster C by:  $\frac{diam(C)}{|C|}$  or  $\frac{avd(C)}{|C|}$ 



## **Evaluating the clustering (Davies-Bouldin index)**

### **Definition (Davies-Bouldin index):**

Let  $k \in \mathbb{N}$  and  $C_1, \dots, C_k$  be the clusters of some data set.

Let  $c_i := cent(C_i)$  for all  $i \in \{1, ..., k\}$ .

Define the **average radius**  $r_i$  of a cluster  $C_i$  as:

$$r_i \coloneqq \frac{1}{|C_i|} \sum_{c \in C_i} d(c_i, c)$$

Define the **Davies-Bouldin index**  $DB(C_1, ..., C_k)$  for  $C_1, ..., C_k$  as:

$$DB(C_1, \dots, C_k) \coloneqq \frac{1}{k} \sum_{i=1}^k \max_{j \neq i} \frac{r_i + r_j}{d(c_i, c_j)}$$

Note: Index is low if average radii are small and clusters are far apart.



## (Naive) Hierarchical clustering algorithm

### Agglomerative (bottom up):

- Initialise each data point as a (singleton) cluster.
   AND initialise representatives for clusters (centroid / clustroid).
- 2. Repeatedly combine two "most similar / closest" clusters to a new one.
  - BY computing in each turn distances between all pairs of clusters.
- 3. Stop w.r.t. certain criterion (e.g. # of clusters, density of clusters, ...)

  BY computing abort parameter each turn.



## (Naive) Hierarchical clustering algorithm

2. Repeatedly combine two "most similar / closest" clusters to a new one.

BY computing in each turn distances between all pairs of clusters.

If we have n data points in total. This part of the algorithm takes  $O(n^3)$  much time.

Conclusion: Hierarchical clustering takes too long for large data sets!



## The (naive) k-means algorithm

- Sometimes also called Lloyd's algorithm, named after Stuart P. Lloyd.
- During the algorithm, the number of clusters is fixed and equal to k.
- Point assignment method, no hierarchical approach.
  - Every data point will be assigned to precisely one cluster.
- Iterative algorithm that assumes Euclidean setting.
- The algorithm alternates between two steps:
  - Assignment step
  - Update step



## The (naive) k-means algorithm

### **Step 0: Initialisation**

- Choose k data points as representatives (centroids), one for each (singleton) cluster  $C_1, \ldots, C_k$ .
  - e.g. all at random, or
  - one at random and each next as far away from former(s) as possible

### **Assignment step:**

(Re)assign each data point p to the cluster  $C_i$  for which the distance between p and the centroid of  $C_i$  is minimal, i.e. to the cluster  $C_i$  that minimises:  $d(p, cent(C_i))$ .



## The (naive) k-means algorithm

### **Assignment step:**

(Re)assign each data point p to the cluster  $C_i$  for which the distance between p and the centroid of  $C_i$  is minimal, i.e. to the cluster  $C_i$  that minimises:  $d(p, cent(C_i))$ .

### **Update step:**

Recalculate the centroids for all clusters.

#### **Termination rule:**

- E.g. run until the process converges / stabilises, or
- stop if only a few / tiny proportion of data points were reassigned.



## (Naive) k-means algorithm: running time

Let  $k \in \mathbb{N}$  be fixed

Let n be the total number of data points we want to cluster.

Let I denote the **number of iteration** the k-means algorithm needs **until convergence** / it stabilises (abort criterion: no changes anymore).

Then the **running time** is: O(knI).

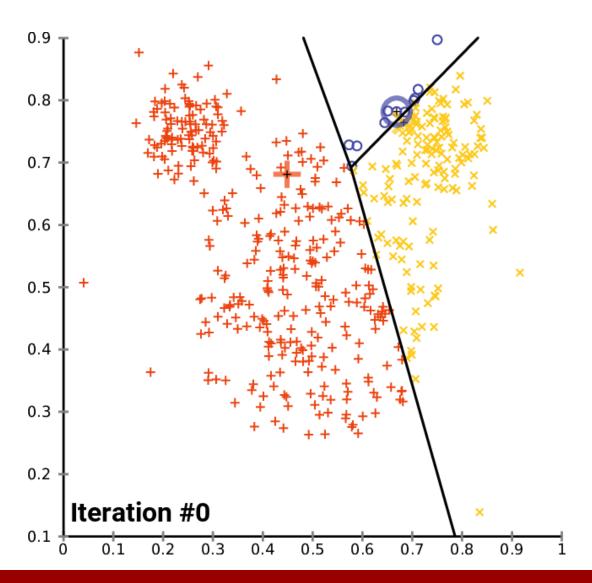
How big is *I*?

Worst case:  $I = 2^{\Omega(\sqrt{n})}$ 

In practice: I can often be considered constant.



## The *k*-means Algorithm



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## What might the right *k* be?

Let  $k \in \mathbb{N}$  and  $C_1, \dots, C_k$  be the clusters of some data set.

Let  $c_i := cent(C_i)$  for all  $i \in \{1, ..., k\}$ .

Define the **average radius**  $r_i$  of a cluster  $C_i$  as:

$$r_i \coloneqq \frac{1}{|C_i|} \sum_{c \in C_i} d(c_i, c)$$

**Often:** Average radii get rapidly smaller while we increase k, up to some value from where on changes are relatively small.

Compute evaluation score: (e.g. Davies-Bouldin index)

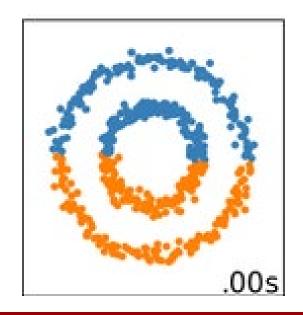


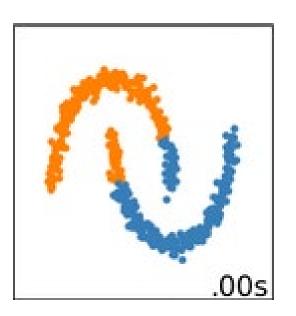
### Problem with *k*-means

In general, k-means (tries) to minimise the following term:

$$\sum_{i=1}^{k} \sum_{c \in C_i} d(c, cent(C_i))$$

Problematic, if (actual) clusters differ a lot in size or are not well-distributed.







## CURE (Clustering Using REpresentatives) algorithm

An algorithm "between" agglomerative and k-means clustering.

**Step 0:** Pick **random sample** of data points fitting into main memory.

**Step 1:** Cluster the sample set, e.g. via agglomerative clustering.

### Step 2:

- Within each cluster  $C_i$ , pick a sample set  $S_i$  of data points which is as "dispersed" as possible.
- Define set of representatives  $R_i$  of  $C_i$  as those points we get by moving each point of  $S_i$  by some fraction towards  $cent(C_i)$ .

**Note:** Points in  $R_i$  are not necessarily actual data points.



**Next:** Merge two clusters if there are two representatives, one from each cluster, whose distance from each other is below some fixed threshold.

→ pick new scattered representatives

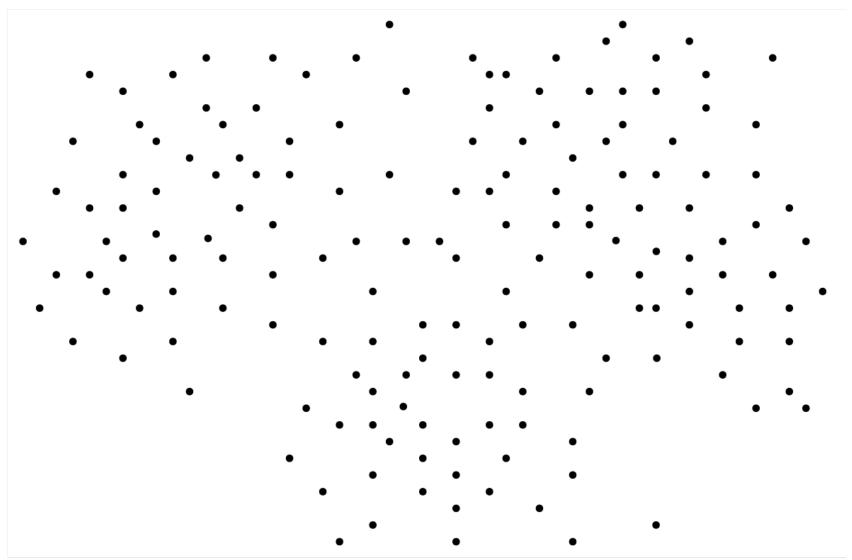
Now start clustering of the whole data by the following rule:

**Assignment rule:** Place data point p into cluster  $C_i$  if p is closest to a representative of  $C_i$  among representatives of all clusters.

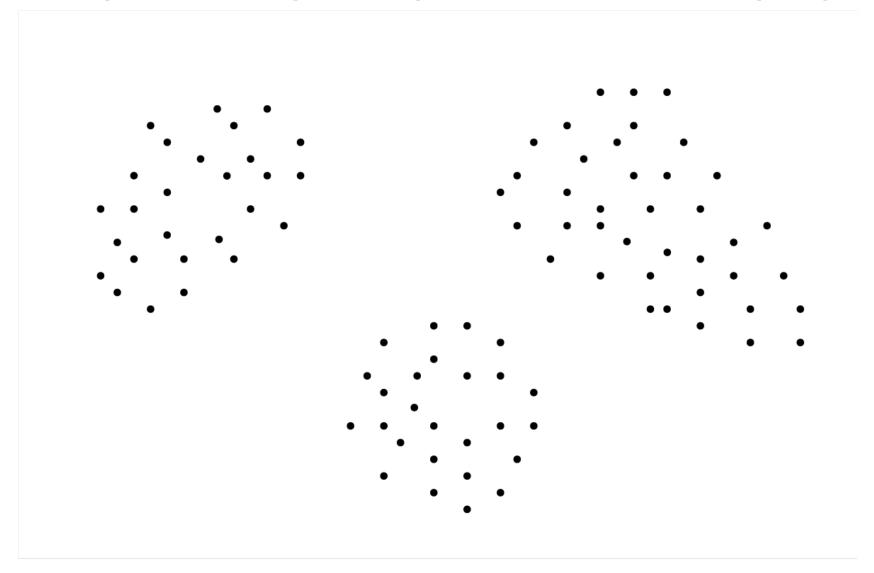
Advantage: Robust for data that is not well-distributed.

Disadvantage: Relatively costly with respect to running time.

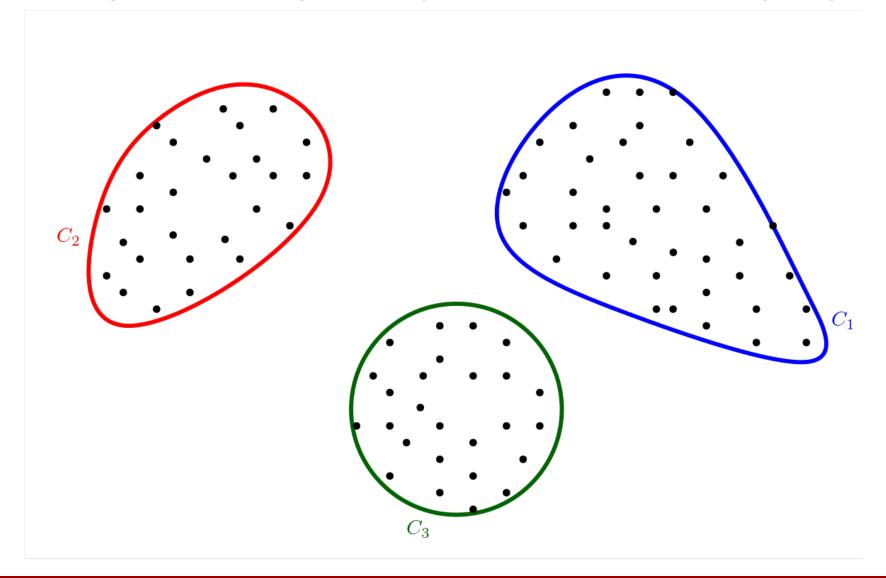




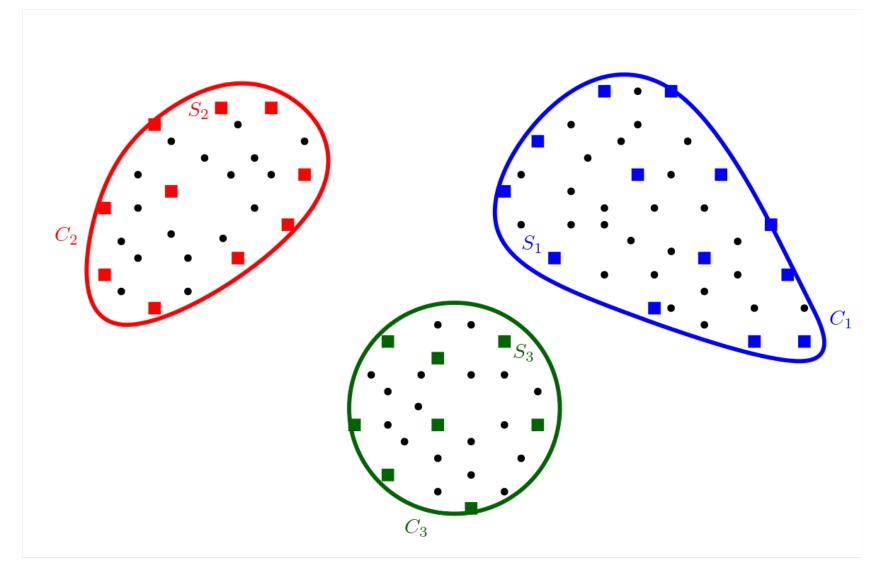




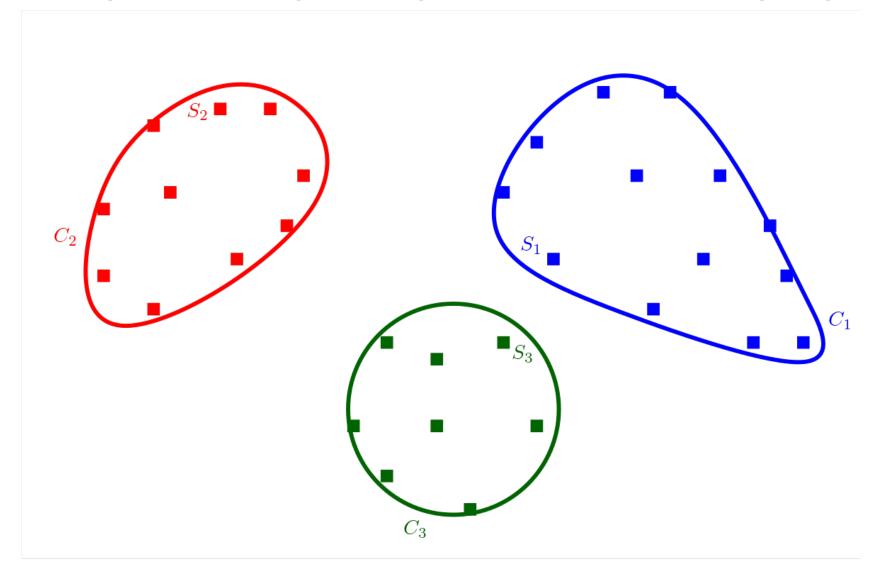




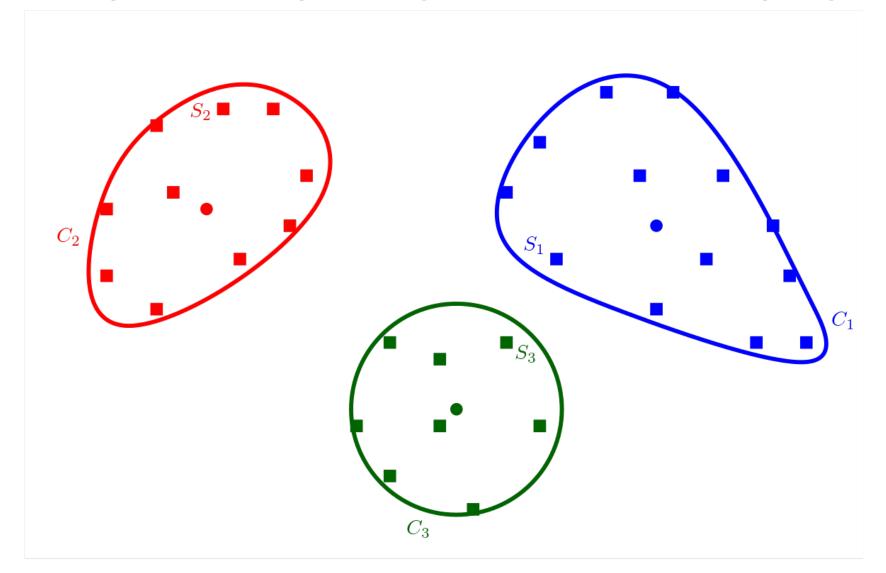




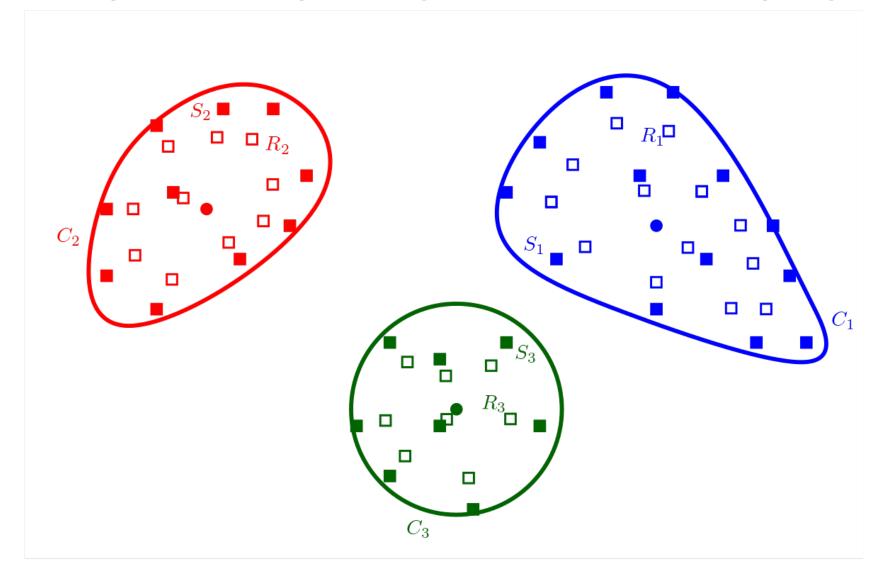




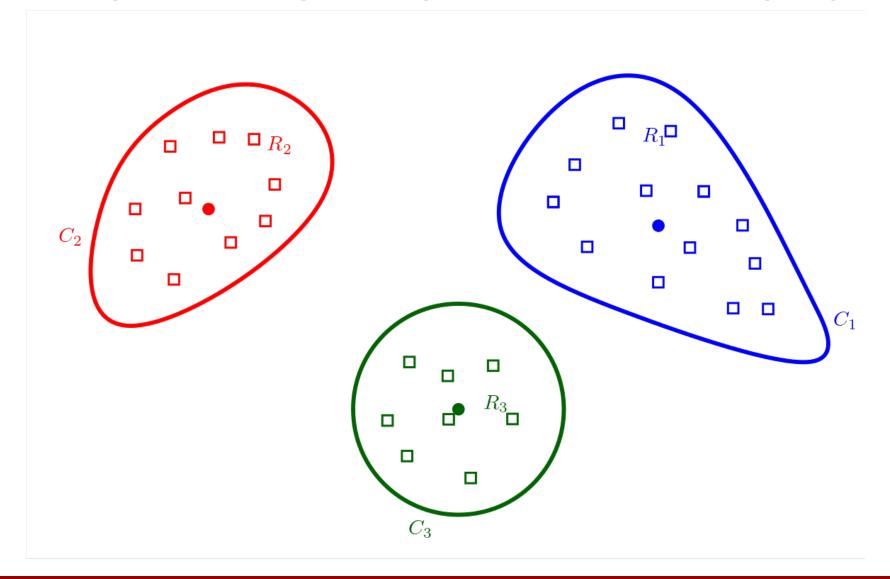




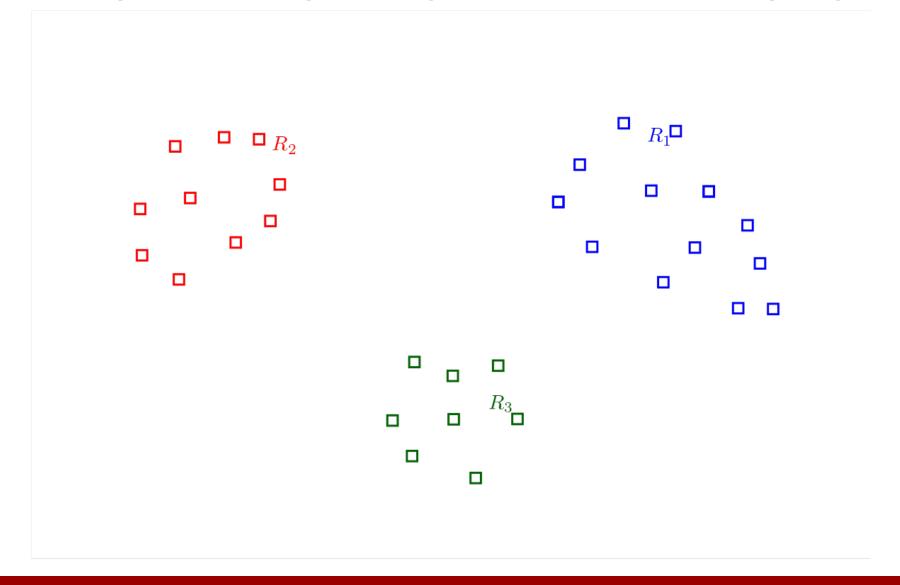




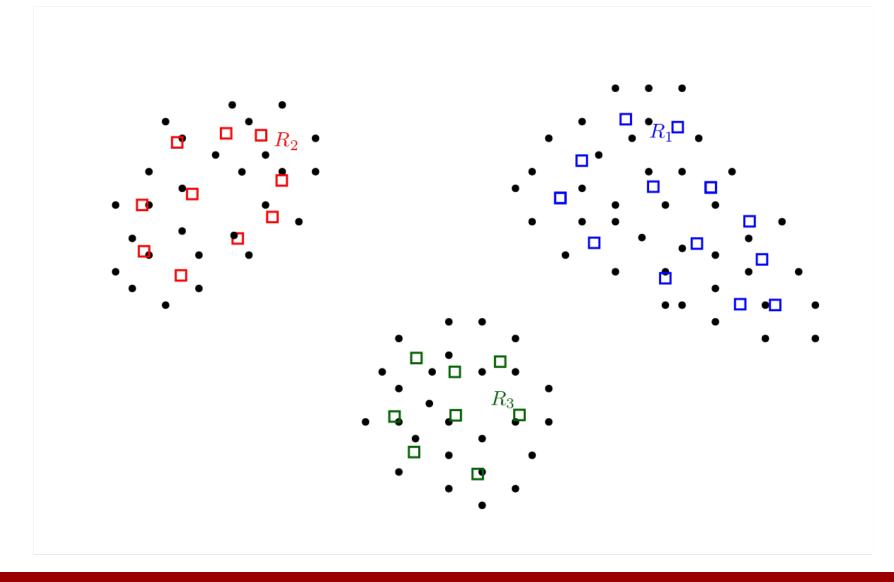




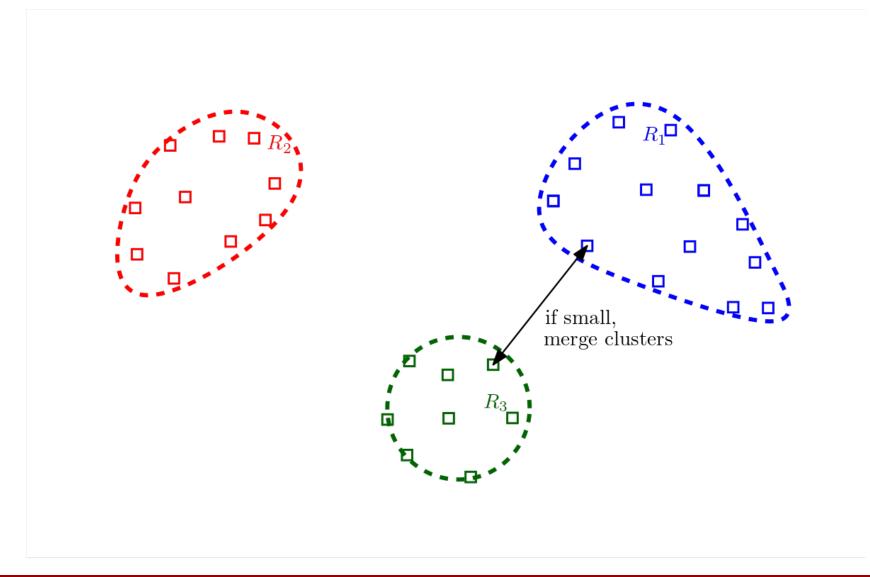




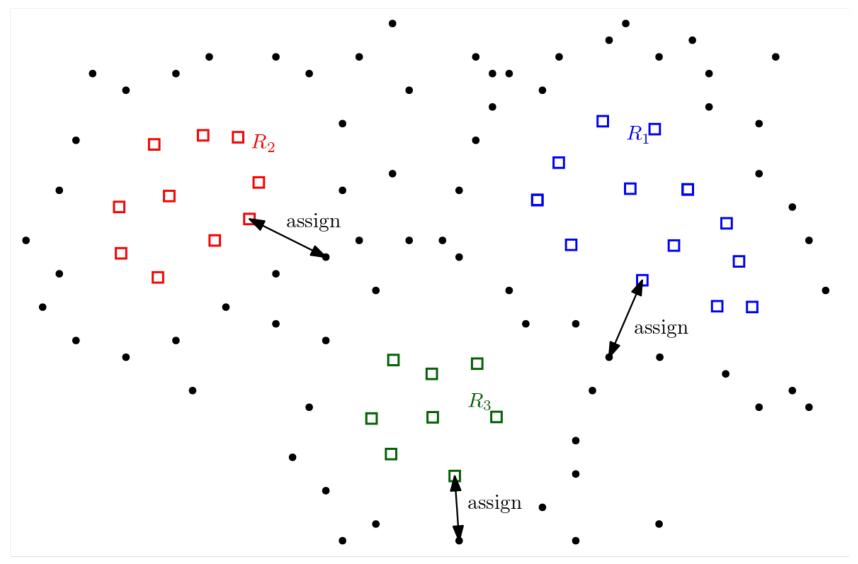




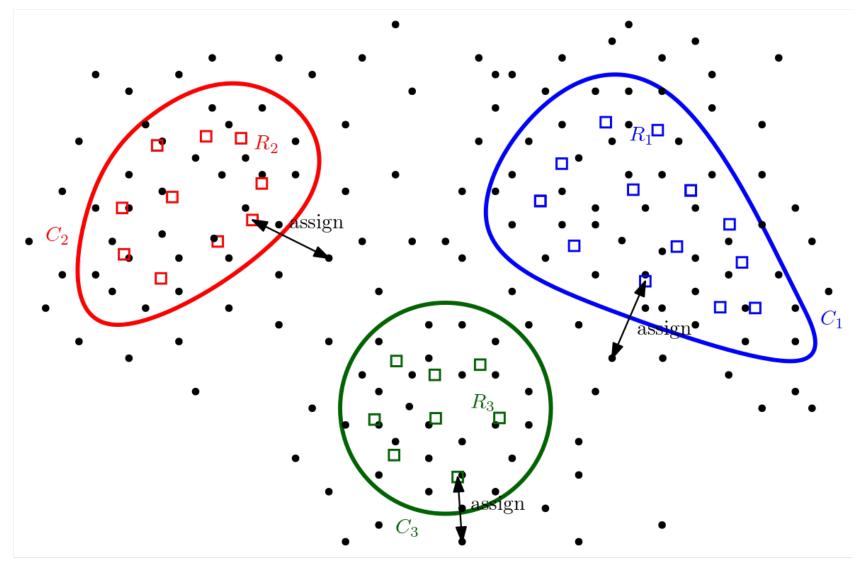














#### Original CURE algorithm

- Perform hierarchical clustering.
- In each step, (re)compute representatives for each cluster.
- Merge two clusters w.r.t. distance between representatives.
- Stop if desired number of clusters is reached (or other abort condition).
- Running time:  $O(n^2 \log(n))$



Name: Density-based spatial clustering of applications with noise

Parameters:  $\varepsilon \in \mathbb{R}, m \in \mathbb{N}$ .

Set *S* of all data points is divided into some different point categories:

•  $p \in S$  is a core point if  $|\{x \in S \mid d(p,x) \le \varepsilon\}| \ge m$ .



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- $x \in S$  is **reachable** from  $p \in S$  if there are points  $x_1, ..., x_k \in S$  such that:
  - $-x_1=p$  and  $x_k=x$ ,
  - $-x_{i+1}$  is directly reachable from  $x_i$  for all  $i \in \{1, ..., k-1\}$ ,
  - $-x_1, \dots, x_{k-1}$  are core points.

A reachable point that is no core point is called a rim point



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A reachable point that is no core point is called a **rim point**.

•  $x \in S$  is an **outlier (or noise) point** if it is not reachable from any  $p \in S$ .



Initialise: Mark each point as unvisited.



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Check each point *p*:

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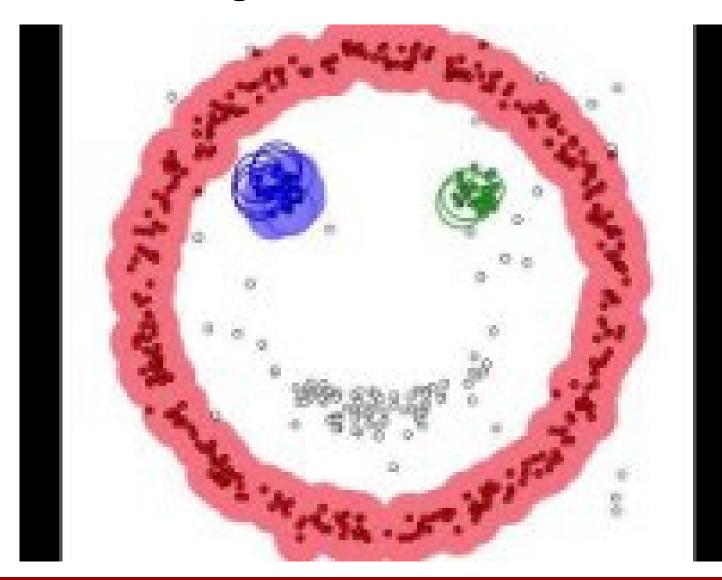
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Running time (worst-case):  $O(n^2)$ 

Advantage: Can identify clusters whose shape is more irregular.





Published by Karthik Ramasamy on Youtube. See this <u>link</u>.