External Memory I

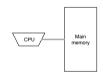
- · Computational Models
- Scanning
- Sorting
- Searching

Philip Bille

External Memory I

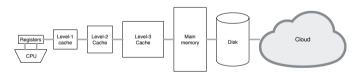
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Computational Models



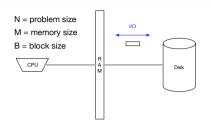
- · (word) RAM Model
 - · Infinite memory of w-bit memory cells
 - Instructions: Memory access, arithmetic operations, boolean operations, control-flow operations, etc.
- · Complexity model.
 - Time = number of instructions.
 - · Space = number of memory cells used.

Computational Models



- · Macbook Pro (late 2023)
 - · CPU: 16 Core M3 Max
- · Registers: ?
- L1 cache: 320 KB per coreL2 cache: 32 MB per core
- · Memory: 128 GB
- · Disk: 1 TB
- Instructions: Memory access, arithmetic operations, boolean operations, control-flow operations, etc.
- · Complexity?

Computational Models



- · I/O model [Aggarwal and Vitter 1988].
 - · Limited memory, Infinite disk
 - Instructions: Disk I/O operations, memory access, arithmetic operations, boolean operations, control-flow operations, etc.
- · Complexity model.
 - · I/Os = Number of disk I/Os
 - · Computation is free (!)

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Scanning

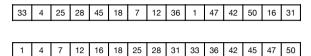


- Scanning. Given an array A of N values (stored in N/B blocks), process all values from left-to
 -right.
- · I/Os. O(N/B).

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Sorting



· Sorting. Given array A of N values (stored in N/B blocks), output the values in increasing order.

Sorting

· Which solutions do we know (on the RAM model)?

External Merge Sort

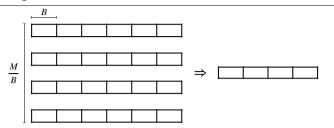
- · Goal. Sorting in O(N/B log_{M/B} (N/B)) I/Os.
- · Solution in 3 steps.
 - · Base case.
 - · External multi-way merge.
 - · External merge sort.

External Merge Sort



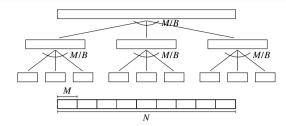
- Base case.
- · Partition N elements into N/M arrays of size M.
- · Load each into memory and sort.
- I/Os. O(N/B)

External Merge Sort



- · Multiway merge algorithm.
 - · Input is N elements in M/B arrays.
 - · Load M/B first blocks into memory and sort.
 - · Output B smallest elements.
 - · Load more blocks into memory if needed.
 - · Repeat.
- I/Os. O(N/B).

External Merge Sort



- · Algorithm.
- · Partition N elements into N/M arrays of size M. Load each into memory and sort.
- · Apply M/B way external multiway merge until left with single sorted array.
- · I/Os
- · Sort N/M arrays: O(N/B) I/Os
- Height of tree O(log_{M/B}(N/M))
- · Cost per level: O(N/B) I/Os.

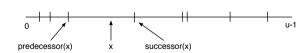
Total I/Os:
$$O\left(\frac{N}{B}\log_{M/B}\frac{N}{M}\right) = O\left(\frac{N}{B}\log_{M/B}\frac{N}{B}\right)$$

External Memory I

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Searching

- · Searching. Maintain a set S ⊆ U = {0, ..., u-1} supporting
- member(x): determine if $x \in S$
- predecessor(x): return largest element in $S \le x$.
- successor(x): return smallest element in $S \ge x$.
- insert(x): set $S = S \cup \{x\}$
- delete(x): set $S = S \{x\}$



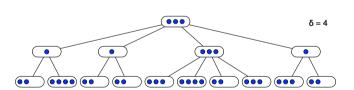
Searching

- · Applications.
 - · Relational data bases.
 - · File systems.

Searching

· Which solutions do we know (on the RAM model)?

B-tree

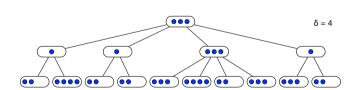


10 42 70

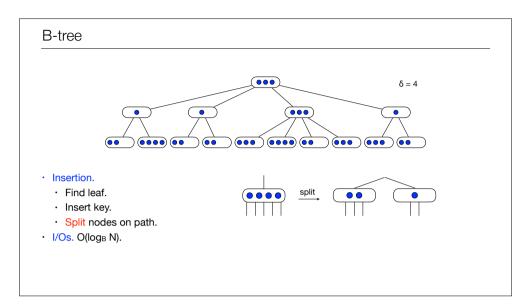
..10 11..42 43..70 71...

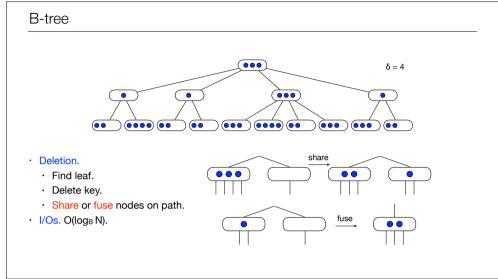
- B-tree of order $\delta = \Theta(B)$ with N keys.
 - · Keys in leaves. Routing elements in internal nodes.
 - Degree between δ/2 and δ.
 - · Root degree between 2 and δ .
 - · Leaves store between $\delta/2$ and δ keys.
 - · All leaves have the same depth.
- Height. $\Theta(\log_{\delta}(N/B)) = \Theta(\log_{B} N)$

B-tree



- · Searching.
- · Find leaf using routing elements.
- · I/Os. O(log_B N).





Basic Bounds

	Internal	External
Scanning	O(N)	scan(N) = O(N/B)
Sorting	O(Nlog N)	$sort(N) = O((N/B)log_{M/B} (N/B))$
Searching	O(log N)	$search(N) = O(log_B(N))$

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- $\cdot \ \text{Sorting}$
- Searching