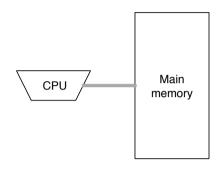
- Computational Models
- Scanning
- Sorting
- Searching

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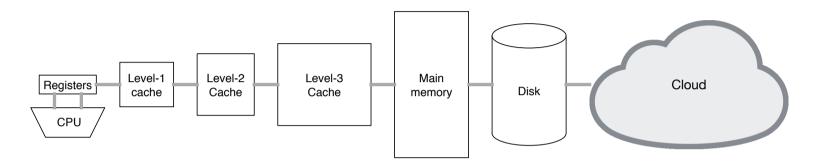
Computational Models



· (word) RAM Model

- · Infinite memory of w-bit memory cells
- Instructions: Memory access, arithmetic operations, boolean operations, control-flow operations, etc.
- · Complexity model.
 - Time = number of instructions.
 - Space = number of memory cells used.

Computational Models



Macbook Pro (late 2023)

· CPU: 16 Core M3 Max

· Registers: ?

· L1 cache: 320 KB per core

· L2 cache: 32 MB per core

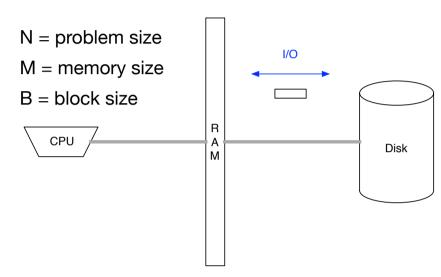
· Memory: 128 GB

· Disk: 1 TB

 Instructions: Memory access, arithmetic operations, boolean operations, control-flow operations, etc.

· Complexity?

Computational Models



- I/O model [Aggarwal and Vitter 1988].
 - · Limited memory, Infinite disk
 - Instructions: Disk I/O operations, memory access, arithmetic operations, boolean operations, control-flow operations, etc.
- · Complexity model.
 - I/Os = Number of disk I/Os
 - Computation is free (!)

- Computational Models
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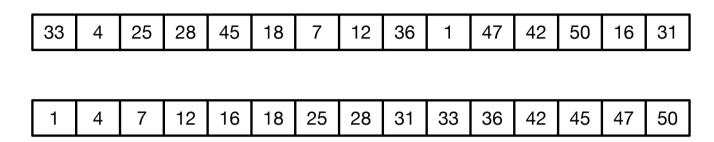
Scanning



- Scanning. Given an array A of N values (stored in N/B blocks), process all values from left-to
 -right.
- I/Os. O(N/B).

- Computational Models
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Sorting

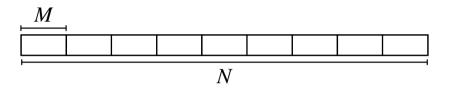


· Sorting. Given array A of N values (stored in N/B blocks), output the values in increasing order.

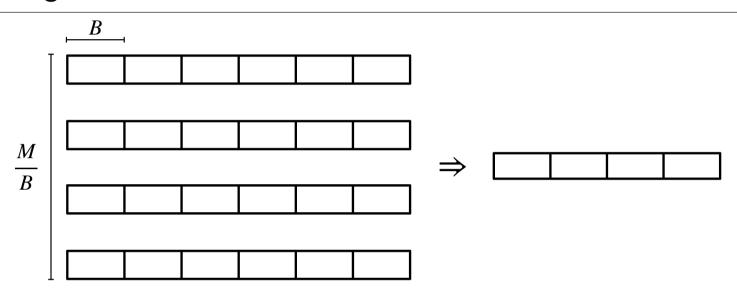
Sorting

Which solutions do we know (on the RAM model)?

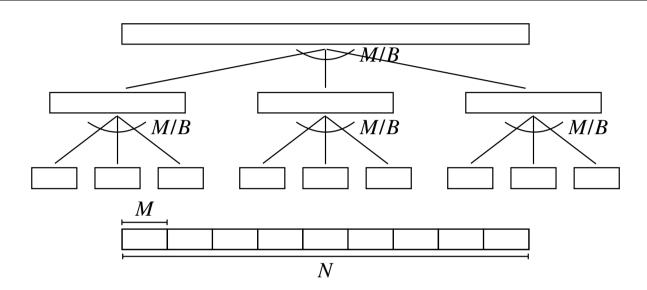
- Goal. Sorting in O(N/B log_{M/B} (N/B)) I/Os.
- Solution in 3 steps.
 - · Base case.
 - External multi-way merge.
 - External merge sort.



- · Base case.
 - Partition N elements into N/M arrays of size M.
 - Load each into memory and sort.
- I/Os. O(N/B)



- Multiway merge algorithm.
 - Input is N elements in M/B arrays.
 - · Load M/B first blocks into memory and sort.
 - Output B smallest elements.
 - · Load more blocks into memory if needed.
 - · Repeat.
- I/Os. O(N/B).



Algorithm.

- Partition N elements into N/M arrays of size M. Load each into memory and sort.
- Apply M/B way external multiway merge until left with single sorted array.

· I/Os.

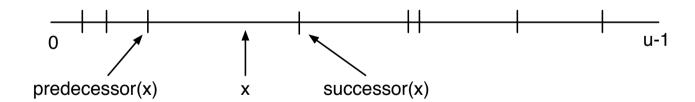
- Sort N/M arrays: O(N/B) I/Os
- Height of tree O(log_{M/B}(N/M))
- Cost per level: O(N/B) I/Os.

Total I/Os:
$$O\left(\frac{N}{B}\log_{M/B}\frac{N}{M}\right) = O\left(\frac{N}{B}\log_{M/B}\frac{N}{B}\right)$$

- Computational Models
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Searching

- Searching. Maintain a set S ⊆ U = {0, ..., u-1} supporting
 - member(x): determine if $x \in S$
 - predecessor(x): return largest element in $S \le x$.
 - successor(x): return smallest element in $S \ge x$.
 - insert(x): set $S = S \cup \{x\}$
 - delete(x): set $S = S \{x\}$

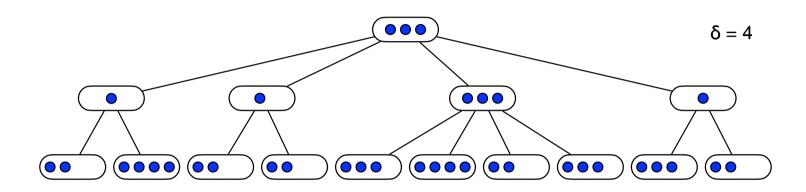


Searching

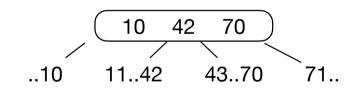
- · Applications.
 - · Relational data bases.
 - · File systems.

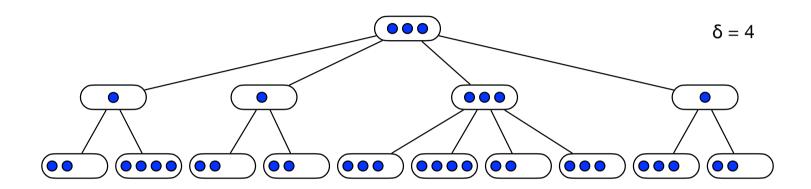
Searching

Which solutions do we know (on the RAM model)?

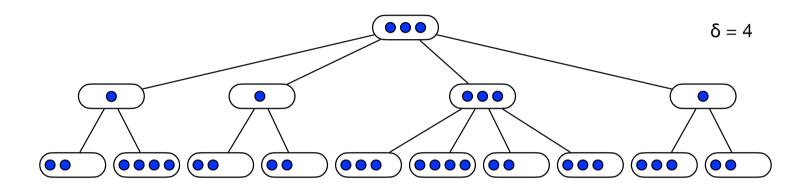


- B-tree of order $\delta = \Theta(B)$ with N keys.
 - Keys in leaves. Routing elements in internal nodes.
 - Degree between $\delta/2$ and δ .
 - Root degree between 2 and δ .
 - Leaves store between $\delta/2$ and δ keys.
 - All leaves have the same depth.
- Height. $\Theta(\log_{\delta}(N/B)) = \Theta(\log_{B}N)$

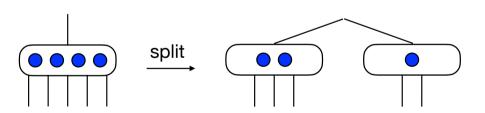


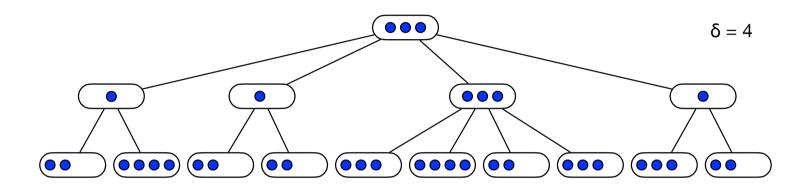


- · Searching.
 - Find leaf using routing elements.
- I/Os. O(log_B N).

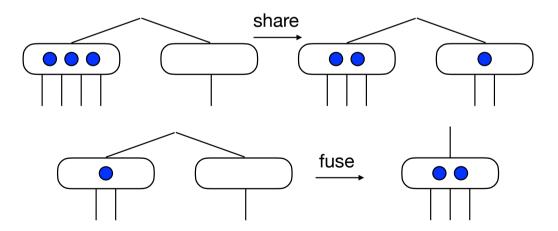


- Insertion.
 - · Find leaf.
 - · Insert key.
 - · Split nodes on path.
- · I/Os. O(log_B N).





- Deletion.
 - · Find leaf.
 - · Delete key.
 - · Share or fuse nodes on path.
- · I/Os. O(log_B N).



Basic Bounds

| | Internal | External |
|-----------|-----------|-------------------------------------|
| Scanning | O(N) | scan(N) = O(N/B) |
| Sorting | O(Nlog N) | $sort(N) = O((N/B)log_{M/B} (N/B))$ |
| Searching | O(log N) | search(N) = $O(log_B(N))$ |

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