Suffix Trees

- String Dictionaries
- Tries
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Inge Li Gørtz

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String Dictionaries

- String dictionary problem. Let S be a string of characters from alphabet Σ. Preprocess S into data structure to support:
- search(P): Return the starting positions of all occurrences of P in S.

- Example.
  - S = yabbadabbado
  - search(abba) = {1,6}

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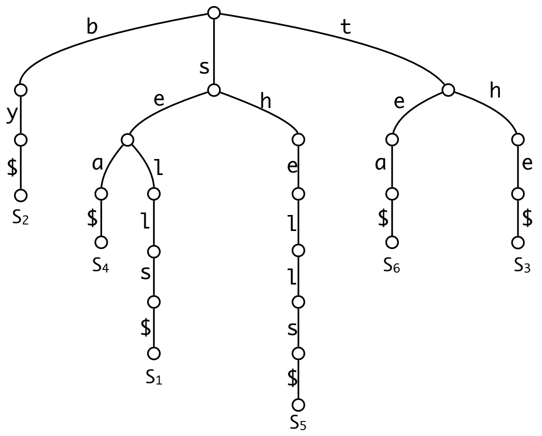
Tries

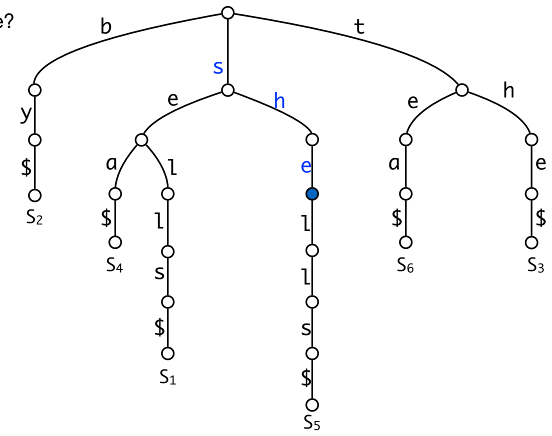
- Text retrieval
- Prefix-free?



- Trie over the strings: sells$, by$, the$, sea$, shells$, tea$, she$.

Tries

- Text retrieval
- Prefix-free?



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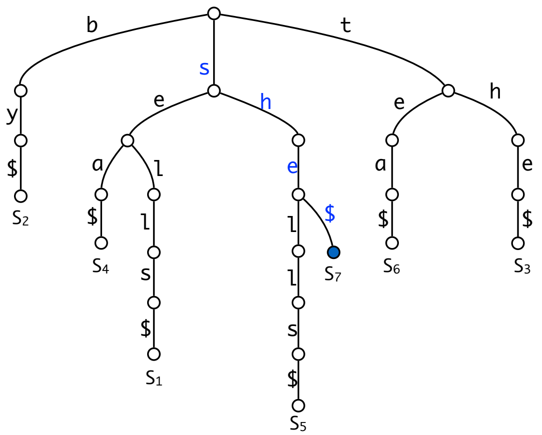


- Trie over the strings: sells$, by$, the$, sea$, shells$, tea$, she$.

Tries

- Properties of the trie. A trie T storing a collection S of s strings of total length *n* from an alphabet of size *d* has the following properties:
    - How many children can a node have?
    - How many leaves does T have?
    - What is the height of T?
    - What is the number of nodes in T?

Trie

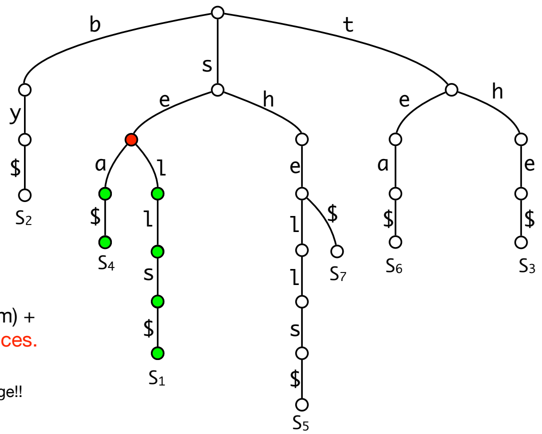
- Properties of the trie. A trie T storing a collection S of s strings of total length *n* from an alphabet of size *d* has the following properties:
  - How many children can a node have? at most d
  - How many leaves does T have? s
  - What is the height of T? length of longest string
  - What is the number of nodes in T? O(n)

Trie

- Search time: O(d) in each node => O(dm).
  - O(m) if d constant.
  - d not constant: use dictionary
    - Perfect hashing O(1)
    - Balanced BST: O(log d)
- Time and space for a trie (for small/constant d):
  - O(m) for searching for a string of length m.
  - O(n) space.
  - Preprocessing: O(n)

Tries

- Prefix search: return all words in the trie starting with "se"



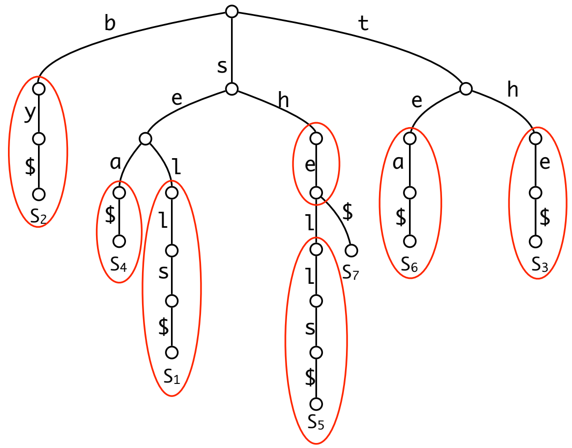
- Time for prefix search: O(m) + time to report all occurrences.

Could be large!!

- Solution: compact tries.

Tries

- Compact trie. Chains of nodes with a single child is merged into a single node.



Tries

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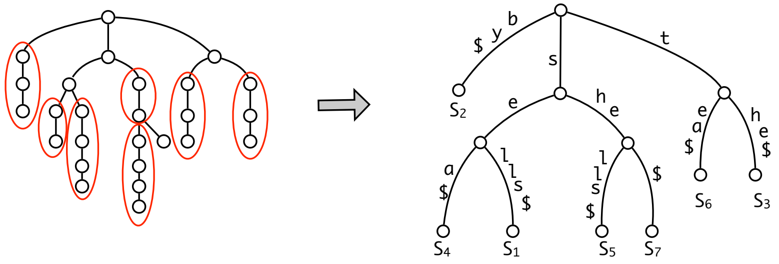

- Properties of the compact trie. A compact trie T storing a collection S of s strings of total length *n* from an alphabet of size *d* has the following properties:
  - Every internal node of T has at least 2 and at most *d* children.
  - T has *s* leaves
  - The number of nodes in T is < 2s.

Trie

- Time and space for a compact trie (constant d).
  - O(m) for searching for a string of length m.
  - O(m + occ) for prefix search, where occ = #occurrences
  - O(s) space.
  - Preprocessing: O(n)

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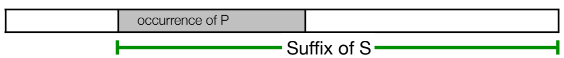
Suffix tree

- String indexing problem. Given a string S of characters from an alphabet Σ. Preprocess S into a data structure to support
  - Search(P): Return starting position of all occurrences of P in S.

- Observation: An occurrence of P is a prefix of a suffix of S.



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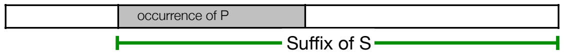
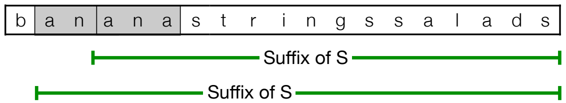
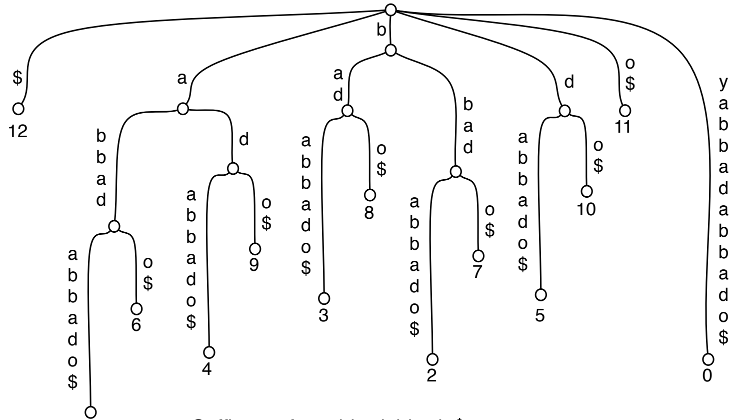


- Example: P = ana.



Suffix Trees

- Suffix trees. The compact trie of all suffixes of S.

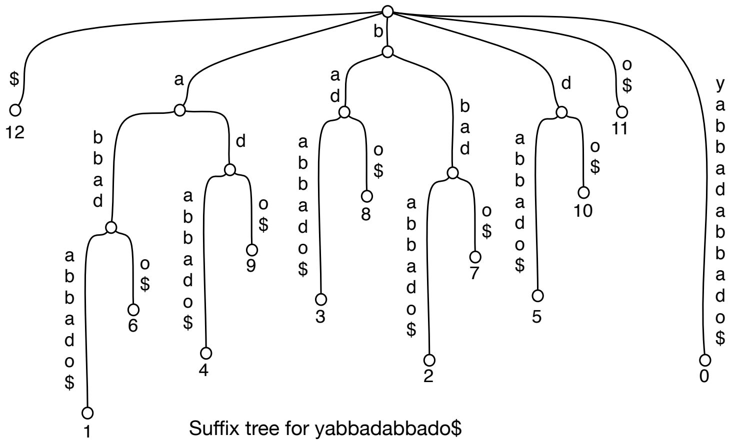


Suffix tree for yabbadabbado$

V

- Suffix trees. The compact trie of all suffixes of S.
- Store S and store node labels by reference to S.

0	1	2	3	4	5	6	7	8	9	10	11	12
y	a	b	b	a	d	a	b	b	a	d	o	$



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Suffix Trees

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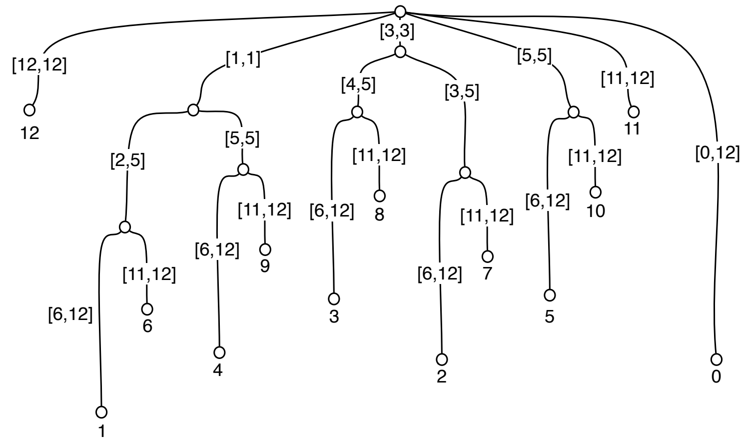
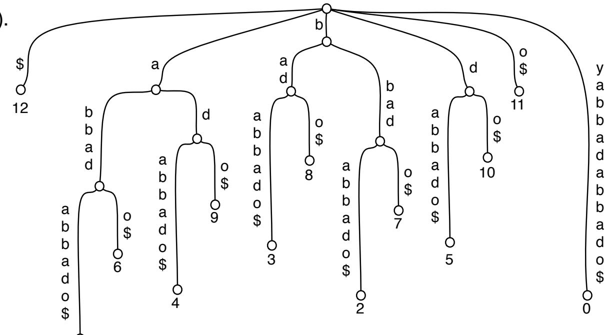


Suffix Trees

- Space.
  - Number of edges + space for edge labels
  - $\implies$ O(n) space
- Preprocessing. O(sort(n,|Σ|))
- sort(n,|Σ|) = time to sort n characters from an alphabet Σ.
- Search(P): O(m+occ).



Suffix tree for yabbadabbado$

Suffix Trees

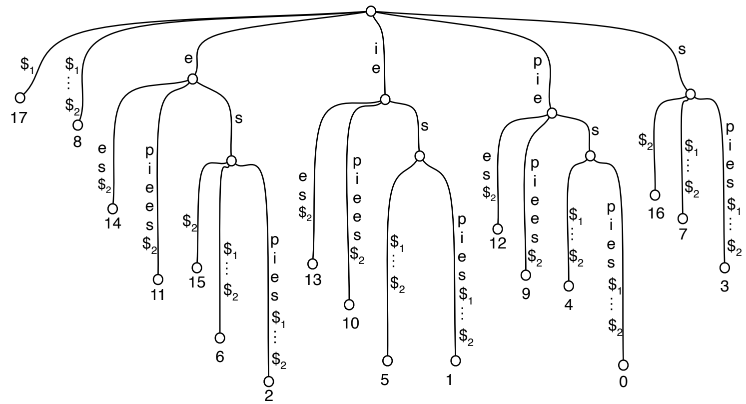
- Theorem. We can solve the string dictionary problem in
  - O(n) space and sort(n,|Σ|) preprocessing time.
  - O(m + occ) time for queries.

Suffix Trees

- Applications.
  - Approximate string matching problems
  - Compression schemes (Lempel-Ziv family, ...)
  - Repetitive string problems (palindromes, tandem repeats, ...)
  - Information retrieval problems (document retrieval, top-k retrieval, ...)
  - ...

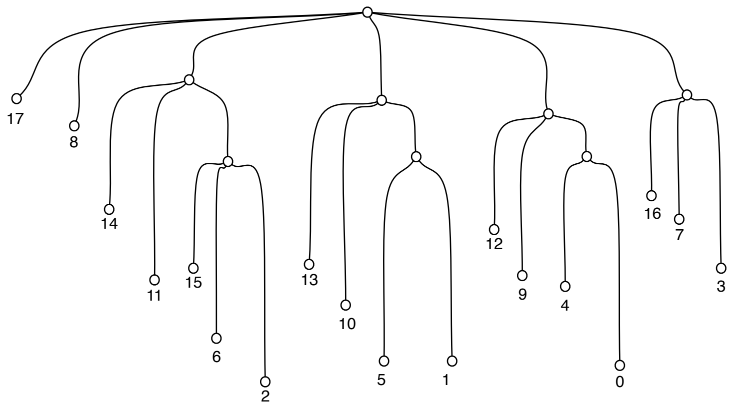
Longest common substring

- Find longest common substring of strings S_1 and S_2.
- Construct the suffix tree over $S_1\$_1S_2\$_2$.
- Example. Find longest common substring of piespies and piepiees:
  - Construct suffix tree of piespies$\$_1$piepiees$\$_2$.



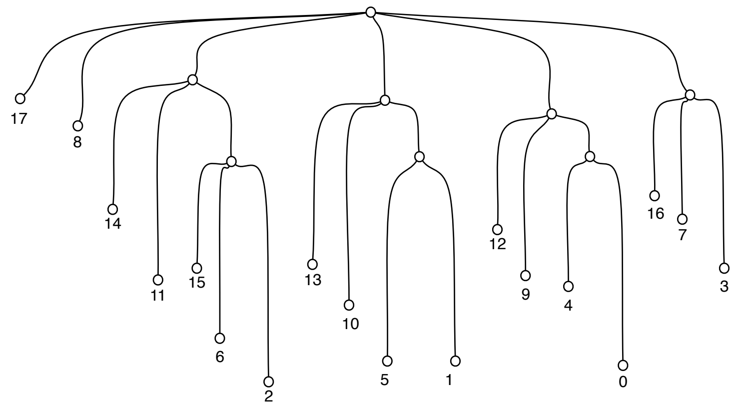
Longest common substring

- Suffix tree of piespies$\$_1$piepiees$\$_2$.



Longest common substring

- Suffix tree of piespies$\$_1$piepiees$\$_2$.
- Mark leaves: ■ = S_1 ▲ = S_2



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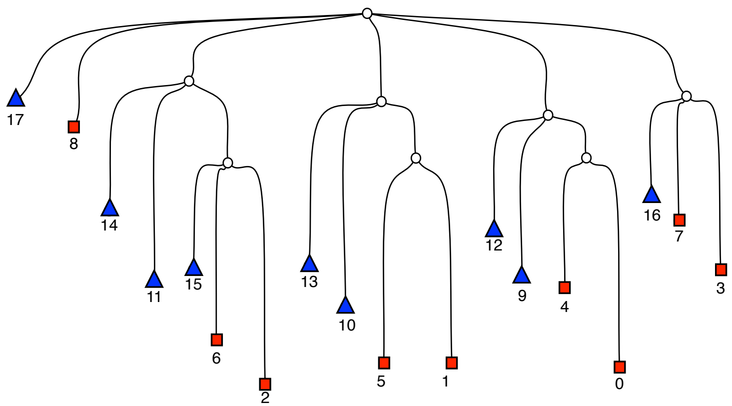

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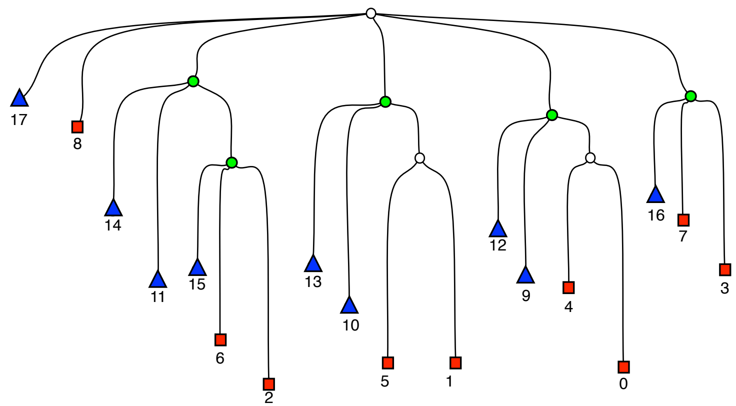
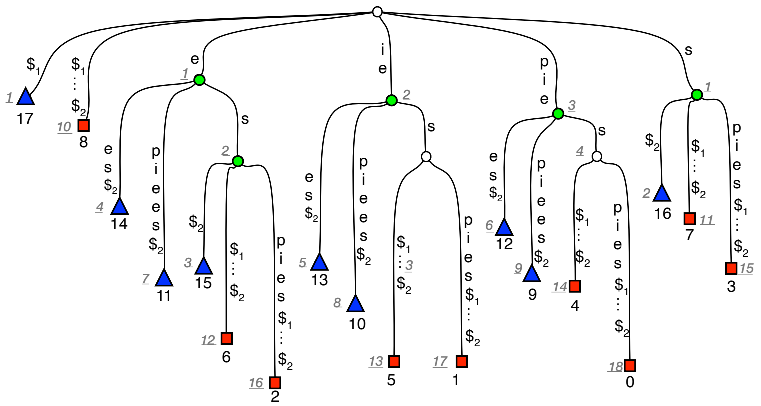


Longest common substring

- Suffix tree of piespies$\$_1$piepiees$\$_2$.
- Mark leaves: ■ = S_1 ▲ = S_2
- Add string depth.



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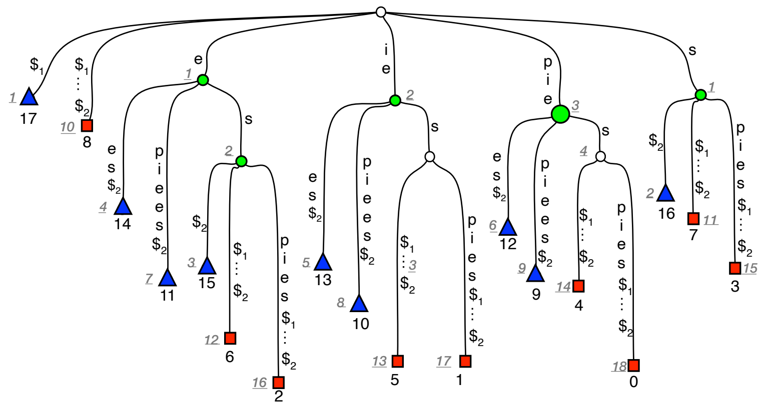

Longest common substring

- Using a suffix tree we can solve the longest common substring problem in linear time (for a constant size alphabets).