Searching and Sorting

- Searching
  - Linear search
  - Binary search
- Sorting
  - Insertion sort
  - Merge sort

Philip Bille

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Searching

- Searching. Given a sorted array A and number x, determine if x appears in the array.

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
1	4	7	12	16	18	25	28	31	33	36	42	45	47	50

Linear Search

- Linear search. Check if each entry matches x.
- Time?
- Challenge. Can we take advantage of the sorted order of the array?

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
1	4	7	12	16	18	25	28	31	33	36	42	45	47	50

Binary Search

- Binary search. Compare x to middle entry m in A.
  - if A[m] = x return true and stop.
  - if A[m] < x continue recursively on the right half.
  - if A[m] > x continue recursively on the left half.
- If array size ≤ 0 return false and stop.

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
1	4	7	12	16	18	25	28	31	33	36	42	45	47	50

Binary Search

```
BINARYSEARCH(A,i,j,x)
   if j < i return false
   m = ⌊i+j/2⌋
   if A[m] = x return true
   elseif A[m] < x return BINARYSEARCH(A,m+1,j,x)
   else return BINARYSEARCH(A,i,m-1,x)     // A[m] > x
```

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
1	4	7	12	16	18	25	28	31	33	36	42	45	47	50

- Time?
- Analysis 1. Analogue of recursive peak algorithm.
  - A recursive call takes constant time.
  - Each recursive call halves the size of the array. We stop when the size is ≤ 0.
  - ⇒ Running time is Θ(log n)

Binary Search

- Analysis 2. Let T(n) be the running time for binary search.
  - Solve the recurrence relation for T(n).

$$T(n) = \begin{cases} T(n/2) + c & \text{if } n > 1 \\ d & \text{otherwise} \end{cases}$$

$$\begin{aligned} T(n) &= T\left(\frac{n}{2}\right) + c \\ &= T\left(\frac{n}{4}\right) + c + c \\ &= T\left(\frac{n}{8}\right) + c + c + c \\ &\;\;\vdots \\ &= T\left(\frac{n}{2^k}\right) + ck \\ &\;\;\vdots \\ &= T\left(\frac{n}{2^{\log_2 n}}\right) + c\log_2 n \\ &= T(1) + c\log_2 n \\ &= d + c\log_2 n \\ &= \Theta(\log n) \end{aligned}$$

Searching

- We can search in
  - Θ(n) time with linear search.
  - Θ(log n) time with binary search.

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Sorting

- Sorting. Given array A[0..n-1] return array B[0..n-1] with same values as A but in sorted order.

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
33	4	25	28	45	18	7	12	36	1	47	42	50	16	31

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
1	4	7	12	16	18	25	28	31	33	36	42	45	47	50

Application

- Obvious.
  - Sort list of names, showGoogle PageRank results, show social media feed in chronological order.
- Non obvious.
  - Data compression, computer graphics, bioinformatics, recommendations systems.
- Easy problem for sorted data.
  - Search, find median, find duplicates, find closest pair, find outliers.

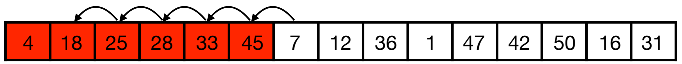
Insertion Sort

- Insertion sort. Start with unsorted array A.
- Proceed left-to-right in n rounds.
- Round i:
  - Subarray A[0..i-1] is sorted.
  - Insert A[i] into A[0..i-1] to make A[0..i] sorted.

4	18	25	28	33	45	7	12	36	1	47	42	50	16	31

Insertion Sort

```
INSERTIONSORT(A, n)
   for i = 1 to n-1
      j = i
      while j > 0 and A[j-1] > A[j]
         swap A[j] og A[j-1]
         j = j - 1
```



- Time?
  - To insert A[i] we use c·i time for constant c.
  - $\Longrightarrow$ total time T(n):

$$T(n) = \sum_{i=1}^{n-1} ci = c\sum_{i=1}^{n-1} i = \frac{cn(n-1)}{2} = \Theta(n^2)$$

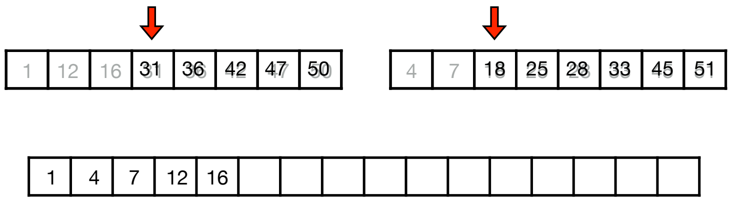
- Challenge. Can we sort faster?

Merge sort

- Merge sort.
  - Idea. Recursive sorting via merging sorted subarray.

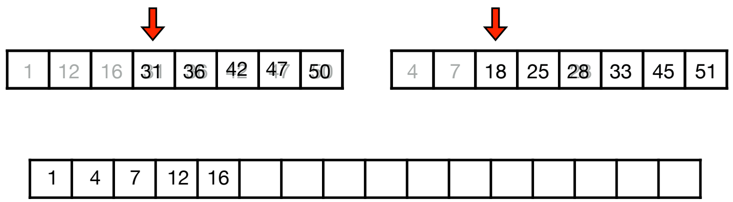
Merge

- Goal. Combine two sorted array into a single sorted array.
- Idea.
  - Scan both arrays left-to-right. In each step:
    - Insert smallest of the two entries in new array.
    - Move forward in array with smallest entry.
    - Repeat until input array exhausted.



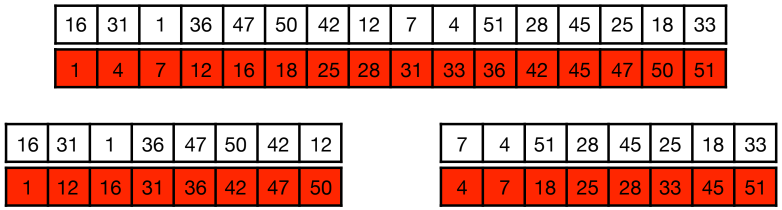
Merge

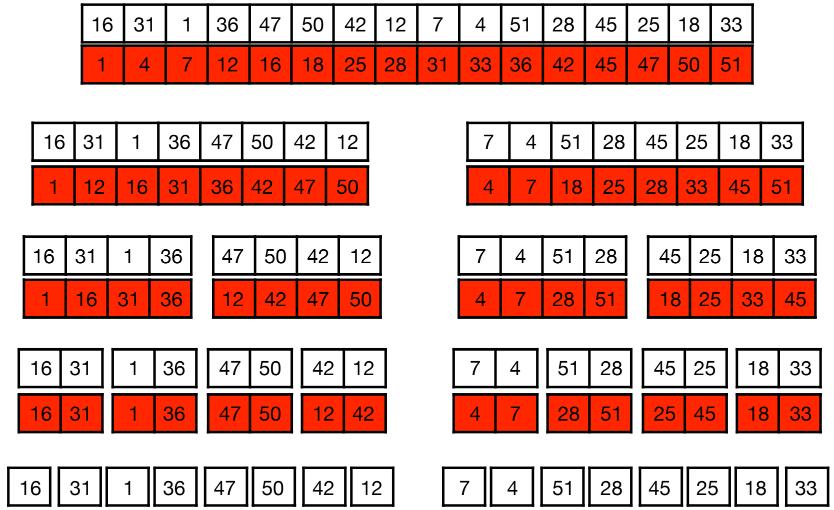
- Time. Merging two arrays A_1 og A_2?
  - Each step take $\Theta(1)$ time.
  - Each step we move forward in one array.
  - $\Rightarrow \Theta(|A_1| + |A_2|)$ time.



Merge sort

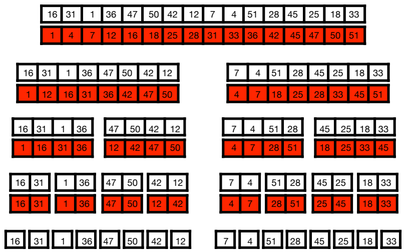
- Merge sort.
- If |A| ≤ 1, return A.
- Otherwise:
  - Split A into halves.
  - Sort each half recursively.
  - Merge the two halves.





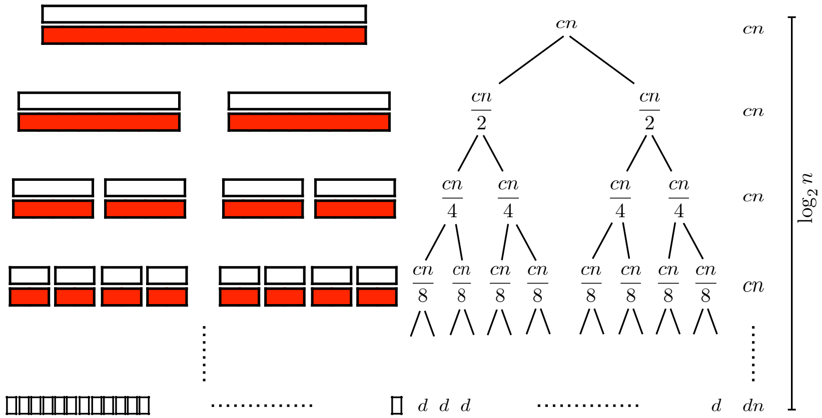
Merge Sort

```
MERGESORT(A,i,j)
   if i < j
      m = ⌊i+j/2⌋
      MERGESORT(A,i,m)
      MERGESORT(A,m+1,j)
      MERGE(A, i, m, j)
```



- Time?
- Construct recursion tree.

Merge Sort



$$T(n) = cn \log_2 n + dn = \Theta(n \log_2 n)$$

Sorting

- We can sort in
  - $\Theta(n^2)$ time with insertion sort.
  - $\Theta(n \log n)$ time with merge sort.

Divide and Conquer

- Merge sort is example of a divide and conquer algorithm.
- Algorithmic design paradigm.
  - Divide. Split problem into subproblems.
  - Conquer. Solve subproblems recursively.
  - Combine. Combine solution for subproblem to a solution for problem.
- Merge sort.
  - Divide. Split array into halves.
  - Conquer. Sort each half.
  - Combine. Merge halves.

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