

# Weekplan: Analysis of Algorithms

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## Reading

*Introduction to Algorithms*, Cormen, Rivest, Leisersons and Stein (CLRS): Chapter 3.

## Exercises

1 [w] **Asymptotic Growth** Arrange the following functions in increasing asymptotic order, i.e., if  $f(n)$  precedes  $g(n)$  then  $f(n) = O(g(n))$ .

$n \log n$     $n^2$     $2^n$     $n^3$     $\sqrt{n}$     $n$

2  **$\Theta$ -notation** Write the following expressions using  $\Theta$ -notation.

$n^2 + n^3/2$   
 $2^n + n^4$   
 $\log_2 n + n\sqrt{n}$   
 $n(n-6)$   
 $4\sqrt{n}$

$8 \log_2^7 n + 34 \log_2 n + \frac{1}{1000} n$   
 $2^{n7} + 5 \log_2^3 n$   
 $n(n^2 - 18) \log_2 n$   
 $n \log_2^4 n + n^2$   
 $n^3 \log_2 n + \sqrt{n} \log_2 n$

3 **Loopy Loops** Analyze the running time of the following loops as a function of  $n$  and express the result in  $\Theta$ -notation.

```
LOOP1(n)
i = 1
while i ≤ n do
  print "*"
  i = 2 · i
end while
```

```
LOOP2(n)
i = 1
while i ≤ n do
  print "*"
  i = 5 · i
end while
```

```
LOOP3(n)
for i = 1 to n do
  j = 1
  while j ≤ n do
    print "*"
    j = 2 · j
  end while
end for
```

4 **Asymptotic Statements** Which of the following statements are true?

$\frac{1}{20}n^2 + 100n^3 = O(n^2)$   
 $\log_2 n + n = O(n)$   
 $2^{\log_2 n} = O(n)$   
 $n^3(n-1)/5 = \Theta(n^3)$   
 $\log_2^2 n + n = \Theta(n)$

$\frac{n^3}{1000} + n + 100 = \Omega(n^2)$   
 $2^n + n^2 = \Omega(n)$   
 $\log_4 n + \log_{16} n = \Theta(\log n)$   
 $n^{1/4} + n^2 = \Theta(n)$   
 $2^{\log_4 n} = \Theta(\sqrt{n})$

**5 Doubling Hypothesis** Solve the following exercises.

- 5.1 [w] The algorithm  $A$  runs in exactly  $7n^3$  time on an input of size  $n$ . How much slower does it run if the input size is doubled?
- 5.2 [BEng] The algorithm  $B$  runs in time respectively 5, 20, 45, 80 and 125 seconds on input of sizes 1000, 2000, 3000, 4000 and 5000. Estimate how long the running time will be of  $B$  on an input of size 6000. What is the running time of  $B$  expressed using  $\Theta$ -notation?
- 5.3 The algorithm  $C$  runs 3 seconds slower each time the size of the input is doubled. What is the running time of  $C$  expressed using  $\Theta$ -notation?

**6 Asymptotic Properties** Solve the following exercises.

- 6.1 CLRS 3.1-1
- 6.2 CLRS 3.1-3
- 6.3 CLRS 3.1-4
- 6.4 [BSc\*] Show that  $\log_2(n!) = \Theta(n \log n)$ . *Hint:* Start by showing the upper bound.
- 6.5 [BSc\*] CLRS 3-2

**7 Generalized Merge Sort** Professor Gørtz suggests the following variant of merge sort called 3-merge sort. 3-merge sort works exactly like normal merge sort except one splits the array into 3 parts instead of 2 that are then recursively sorted and merged. Solve the following exercises.

- 7.1 Show it is possible to merge 3 sorted arrays in linear time.
- 7.2 Analyze the running time of 3-merge sort.
- 7.3 [\*] Generalize the algorithm and the analysis of 3-merge sort to  $k$ -merge sort for  $k > 3$ . Is  $k$ -merge sort an improvement over the standard 2-merge sort?

**8 Maximal Subarray** Let  $A[0..n-1]$  be an array of integers (both positive and negative). A *maximal subarray* of  $A$  is a subarray  $A[i..j]$  such that the sum  $A[i] + A[i+1] + \dots + A[j]$  is maximal among all subarrays of  $A$ . Solve the following exercises.

- 8.1 [w] Give an algorithm that finds a maximal subarray of  $A$  in  $O(n^3)$  time.
- 8.2 Give an algorithm that finds a maximal subarray of  $A$  in  $O(n^2)$  time. *Hint:* Show it is possible to compute the sum of any subarray in  $O(1)$  time.
- 8.3 [\*\*] Give a divide and conquer algorithm that finds a maximal subarray of  $A$  in  $O(n \log n)$  time.
- 8.4 [\*\*] Give an algorithm that finds a maximal subarray of  $A$  in  $O(n)$  time.